

25 Nov 2024

# Max Cut and SDP

Obs. If  $G = (V, E)$  is undir graph with  $w(u, v) = \begin{cases} 1 & \text{if } \{u, v\} \in E \\ 0 & \text{o.w.} \end{cases}$   
 and  $x_u \in \{\pm 1\}^V$  then

$$\langle x, L_G x \rangle = \sum_{\{u, v\} \in E} (x_u - x_v)^2$$

$$= 4 \cdot \text{Cut}(x)$$

where  $\text{Cut}(x) = \#$  of edges labeled with  $\{+1, -1\}$  by  $x$ .

So  $\text{Max Cut}(G) = \max \{ \langle x, L_G x \rangle \mid x \in \{\pm 1\}^V \}$ .

So... apply random thresholding to the eigenvector  $x$  corresponding to  $\lambda_{\max}(L_G)$ ?

E.g.  $G = K_n$  clique on  $n$  vertices

$$L_G = \begin{bmatrix} n-1 & -1 & -1 & \dots & -1 \\ -1 & n-1 & & & \\ -1 & & \ddots & & \\ \vdots & & & \ddots & \\ -1 & & & & n-1 \end{bmatrix} = n \cdot \text{Id}_n - \mathbf{1}\mathbf{1}^T$$

Make  $w(u, v) = \begin{cases} 2 & \text{if } u=v_1 \text{ or } v=v_1 \\ 1 & \text{otherwise} \end{cases}$

For this  $G$

$$L_G = \begin{bmatrix} 2(n-1) & -2 & & & -2 \\ -2 & n & & & \\ \vdots & & \ddots & & \\ -2 & -1 & & & n \end{bmatrix} = (n+1)I_n - \mathbf{1}\mathbf{1}^T + ye_1^T + e_1y^T$$

where  $y^T = \left[ \frac{n-2}{2} \quad -1 \quad -1 \quad \dots \quad -1 \right]$ .

Now the  $\lambda_{\max}$  eigenvector is nearly parallel to  $e_1$  and very misaligned with genuine max-cut vectors.

Observation.  $\langle x, L_G x \rangle = x^T L_G x$   
 $= \text{Tr}(L_G x x^T)$ .

So  $\max \left\{ \langle x, L_G x \rangle \mid \langle x, x \rangle = n \right\}$

is the same as:

$$\max \text{Tr}(L_G X)$$

$$\text{s.t. } X \in \mathbb{R}^{n \times n}, \quad X = X^T$$

$$\text{Tr}(X) = n$$

$$X \succeq 0$$

(i.e.  $X$  is PSD)

$$\text{rank}(X) = 1$$

Instead try to solve:

(GW-SDP)    max     $\text{Tr}(L_G \cdot X)$   
                  s.t.     $X \in \mathbb{R}^{n \times n}$ , symmetric, PSD  
                           $\text{Tr}(X) = n$  [redundant]  
                           $x_{ii} = 1 \quad \forall i \in [n]$

} semidef. program.

~~$\text{rank}(X) = 1$~~

SDP: Maximize/minimize a linear function of the entries of a square matrix subj. to:

- the matrix is symmetric PSD
- any number of additional linear inequality constraints on the matrix entries.

Solve GW-SDP.

Let  $X = WW^T$ .

Let  $\{w_1, \dots, w_n\}$  be columns of  $W$ .

$X_{ij} = \langle w_i, w_j \rangle$ .

Sample unit vector  $u \in \mathbb{R}^n$   
uniformly at random.

$$S = \left\{ i \mid \langle w_i, u \rangle > 0 \right\}$$

$$V - S = \left\{ i \mid \langle w_i, u \rangle \leq 0 \right\}.$$

Why does this work?