

20 Oct 2023

Simplex Algorithm and Strong Duality

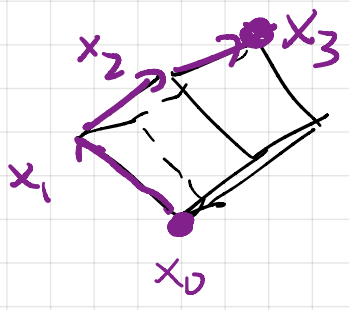
$$\begin{aligned} \max \quad & c^T x \\ \text{s.t.} \quad & Ax \preceq b \\ & x \succeq 0 \end{aligned}$$

$$\begin{aligned} \min \quad & b^T y \\ \text{s.t.} \quad & A^T y \preceq c \\ & y \succeq 0 \end{aligned}$$

The solution set of $\{Ax \preceq b, x \succeq 0\}$ is a polyhedron contained in the positive orthant, $\mathbb{R}_{\geq 0}^n$.

Vertices of the polyhedron correspond to n -tuples of (linearly independent) tight constraints.

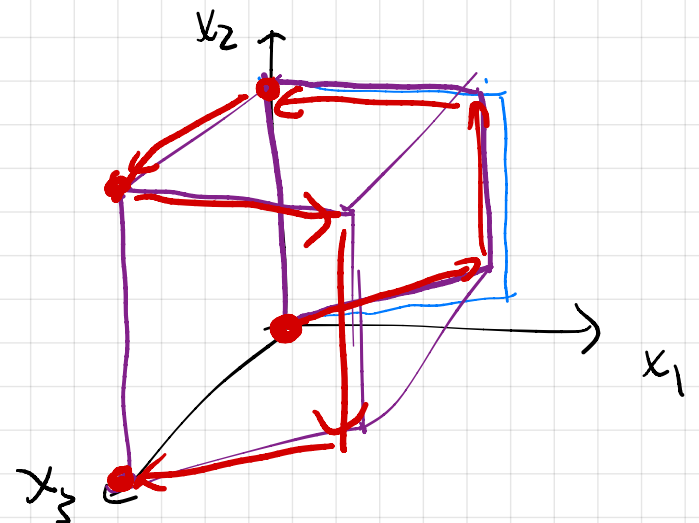
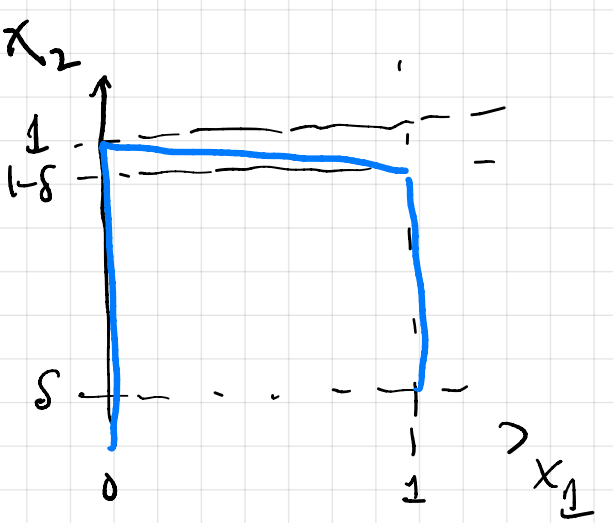
Simplex method is a local search algorithm that starts at any vertex and walks to a neighboring vertex where the objective function is (weakly) greater, until it cannot find any such neighbor.



A polyhedron in n dimensions defined by m constraints (plus non-negativity inequalities) has $\leq \binom{m+n}{n}$ vertices.

Klee-Minty Cubes:

$$\begin{aligned} \max \quad & x_n \\ \text{s.t.} \quad & 0 \leq x_i \leq 1 \\ & \delta x_i \leq x_{i+1} \leq 1 - \delta x_i \quad \forall i < n. \end{aligned}$$



Equational form:

$$\begin{aligned} \max \quad & c^T x \\ \text{s.t.} \quad & Ax \preceq b \\ & x \succeq 0 \end{aligned}$$

$$\begin{aligned} \max \quad & c^T x \\ \text{s.t.} \quad & Ax + w = b \\ & x, w \succeq 0 \end{aligned}$$

vector of "slack variables"

$$\begin{aligned} \max \quad & 2x_1 + 3x_2 \\ \text{s.t.} \quad & x_1 + x_2 \leq 8 \\ & 2x_1 + x_2 \leq 12 \\ & x_1 + 2x_2 \leq 14 \\ & x_1, x_2 \geq 0 \end{aligned}$$

$$\begin{bmatrix} 1 & 1 \\ 2 & 1 \\ 1 & 2 \end{bmatrix} \begin{bmatrix} x_1 \\ x_2 \end{bmatrix} \preceq \begin{bmatrix} 8 \\ 12 \\ 14 \end{bmatrix}$$

$$\begin{aligned} \max \quad & 2x_1 + 3x_2 \\ \text{s.t.} \quad & x_1 + x_2 + w_1 = 8 \\ & 2x_1 + x_2 + w_2 = 12 \\ & x_1 + 2x_2 + w_3 = 14 \\ & \vec{x}, \vec{w} \succeq 0 \end{aligned}$$

$$\begin{bmatrix} 1 & 1 & 1 & 0 & 0 \\ 2 & 1 & 0 & 1 & 0 \\ 1 & 2 & 0 & 0 & 1 \end{bmatrix} \begin{bmatrix} x_1 \\ x_2 \\ w_1 \\ w_2 \\ w_3 \end{bmatrix} = \vec{b}$$

Suppose we start from $x_2 = w_2 = 0$.

