

4/23: CBFs: Theory & Applications

Announcements:

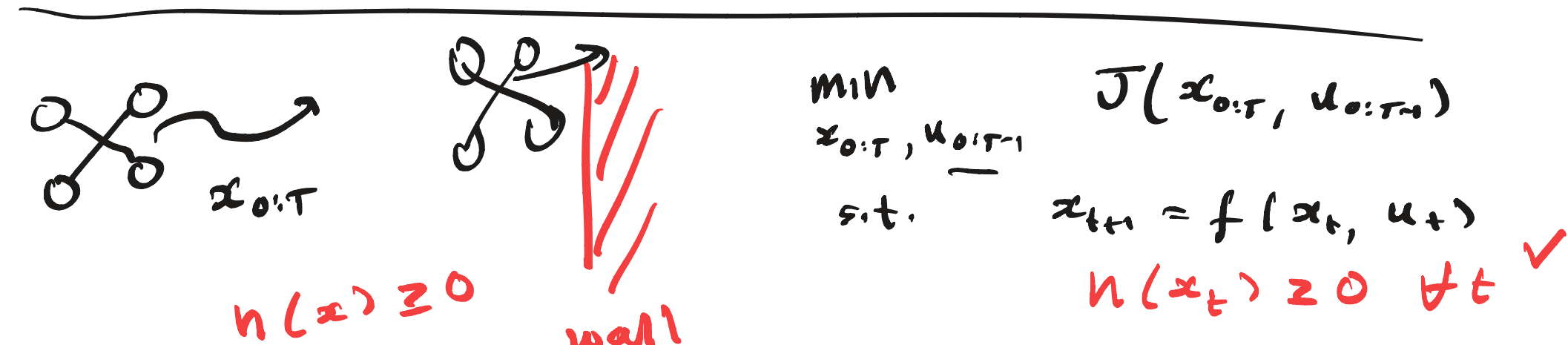
- HW4 due today (4/23)
- project videos due 5/4 (mon) (no late days!)
- project report due 5/11 (mon)
  - ↳ 8 pg. max, results + "lessons learned"
- guest lecture: Prof. Cauligi, 4/30

Last time:

- MBAL

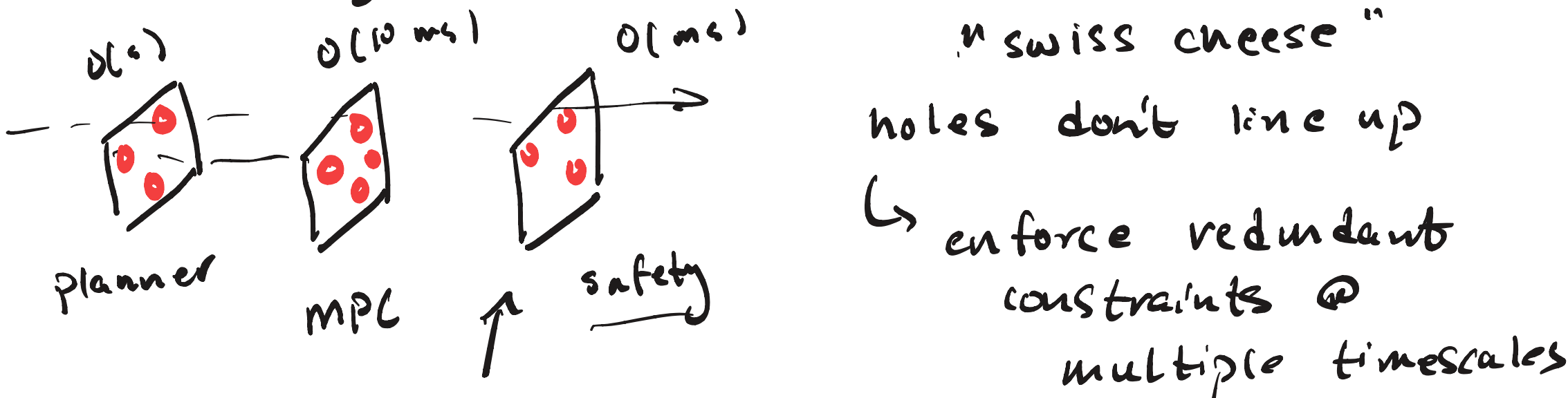
Today:

- CBFs!



Issues: 1) horizon ("recursive feas.")  $u_t \in \mathcal{U}$   
 2) modeling error / mismatch

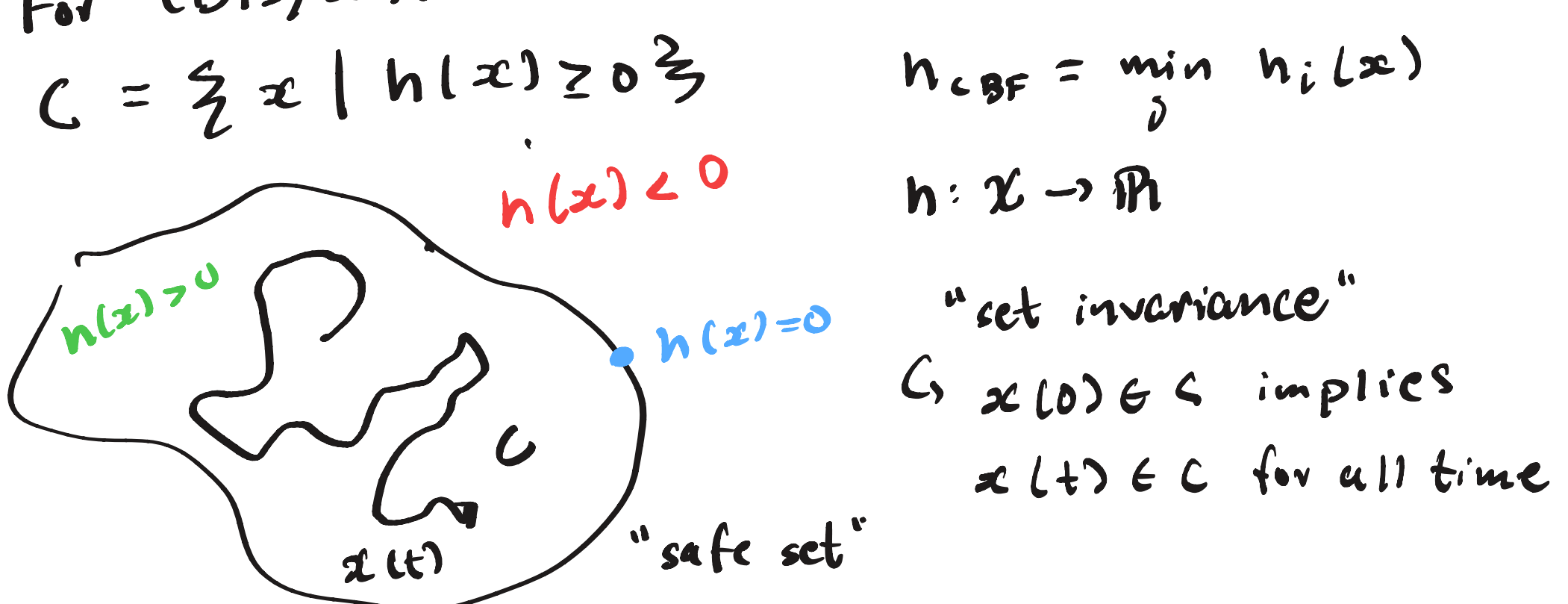
Q: How do you actually get a robot to be safe?



1. Theory of CBFs
2. Applications

In this class "safety" looks like  $n(x) \geq 0$

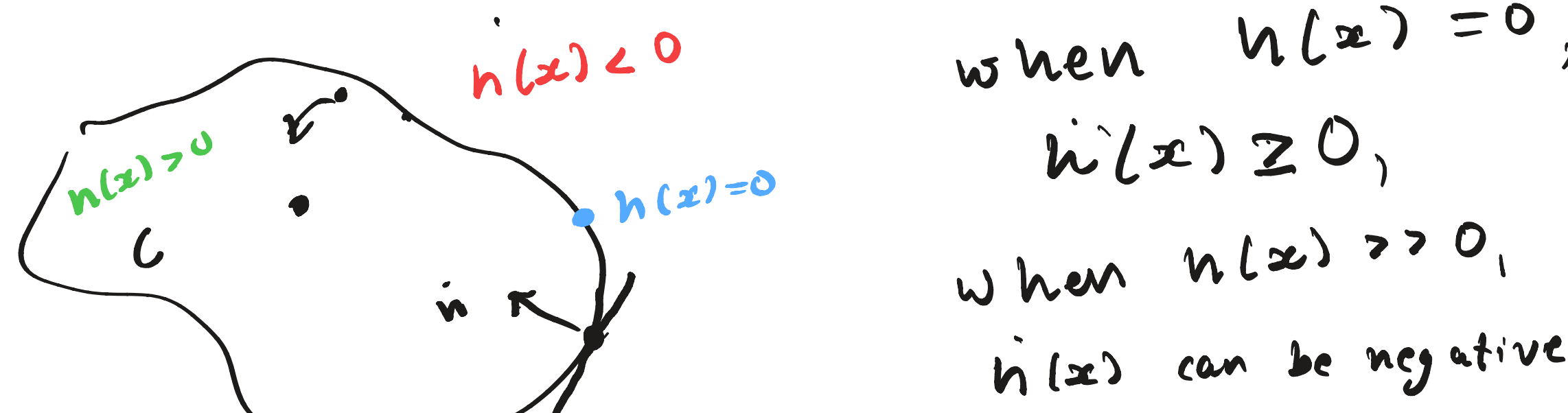
For CBFs/controls  $C = \{x \mid n(x) \geq 0\}$



Key insight: Because  $n(x)$  is one number that determines safety, we can just think about  $\dot{n}(t)$ .

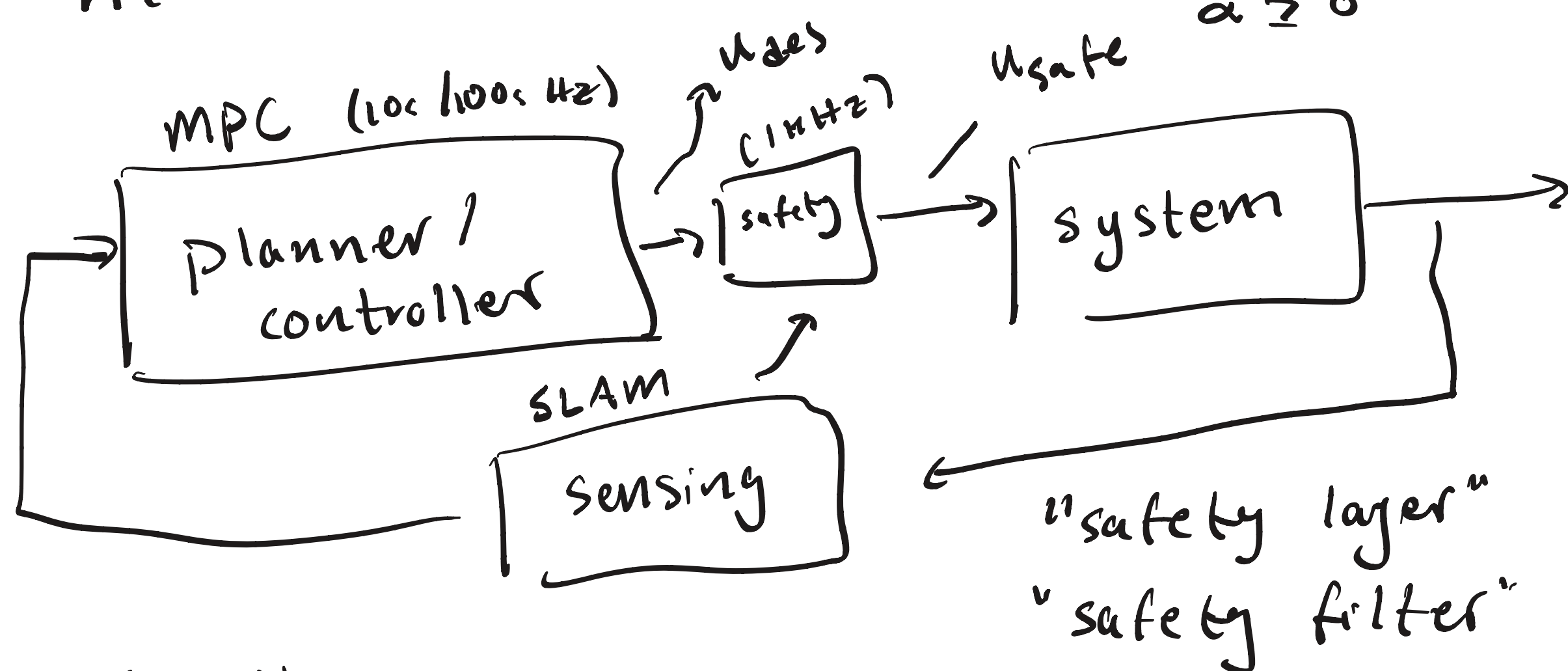
- $\dot{n}(t) > 0 \Rightarrow$  becoming safer
- $\dot{n}(t) < 0 \Rightarrow$  becoming less safe

Q: For safety, what should  $\dot{n}(t)$  look like?



$$\dot{n}(t) = -\alpha n(t)$$

$$\dot{n}(x) \geq -\alpha n(x) \quad \alpha \geq 0$$



Key idea: we want something fast that minimally changes  $u_{des}$  to enforce safety

Safety layer:

$$\min_u \|u - u_{des}\|^2 \quad \text{s.t.} \quad \dot{n}(x, u) \geq -\alpha n(x)$$

$$\dot{n}(x, u) = \frac{dh(x(t))}{dt} = \frac{\partial h^T}{\partial x} \dot{x}$$

Common form  $\dot{x} = f(x) + g(x)u$

$$\dot{n}(x, u) = \frac{\partial h^T}{\partial x} (f(x) + g(x)u)$$

$$\frac{\partial h^T}{\partial x} f(x) + \frac{\partial h^T}{\partial x} g(x)u \geq -\alpha n(x) \quad \text{"active set"}$$

$$\min_u \|u_{des} - u\|^2 \quad \text{s.t.} \quad a^T u \leq b \quad u^* = \begin{cases} u_{des} & a^T u_{des} \leq b \\ \Pi(u_{des}) & \text{otherwise} \end{cases}$$

Discrete time

$$n(x_{k+1}) \geq \alpha n(x_k), \quad \alpha \in [0, 1]$$

$$\min_{u_k} \|u_k - u_{des}\|^2 \quad \text{s.t.} \quad n(f(x_k, u_k)) \geq \alpha n(x_k)$$

Practical details:

- i) Recursive feasibility
- ii) modeling / state est. error
- iii)  $n(x)$  is an open design problem