

21 Feb

Multiplying polynomials and convolution.

Example . $(2x + 1)(4x - 3) = 8x^2 - 2x - 3$

$$\begin{array}{r} 2x + 1 \\ 4x - 3 \\ \hline -6x - 3 \\ 8x^2 + 4x \\ \hline 8x^2 - 2x - 3 \end{array}$$

Quadratic time
 $O(\text{degree}^2)$
arithmetic operations

If $A(x) = A_1x + A_0$

$$B(x) = B_1x + B_0$$

and $C(x) = A(x) B(x)$

then $C(x) = A_1 B_1 x^2 + [A_0 B_1 + A_1 B_0]x + A_0 B_0$
 $= P_2 x^2 + (P_1 - P_2 - P_0)x + P_0$

where $P_0 = A_0 B_0$

$$P_1 = A_1 B_1$$

$$P_2 = (A_0 + A_1)(B_0 + B_1)$$

This is the basis of an $O(n^{\log 3})$ algorithm

for multiplying degree n polynomials,
 similar to Karatsuba's.

Why multiplying polynomials matters..

1. Basic operation in algebra.
2. Underpinning of integer multiplication (crypto)

A binary number of n bits is just a degree $n-1$ polynomial evaluated at $x=2$, with $\{0,1\}$ coefficients.

3. Polynomials can represent signals (sequences of numbers).

Multiplication represents convolution.

Ex. Take the sequence
 $\dots, 0, 0, 4, 8, 2, 0, 0, \dots$

and replace every element with the average of the preceding and following ones.

$$\begin{array}{r} (4x^3 + 8x^2 + 2x + 0) \\ \times \frac{1}{2}x^2 + \frac{1}{2} \\ \hline 2x^5 + 4x^4 + x^3 + 0 \\ 2x^5 + 4x^4 + x^3 + 0 \\ \hline 2x^5 + 4x^4 + 3x^3 + 4x^2 + x + 0 \\ 2 \quad 4 \quad 3 \quad 4 \quad 1 \quad 0 \end{array}$$

$$\text{If } a_0, a_1, a_2, \dots, a_n$$

$$b_0, b_1, b_2, \dots, b_m$$

are two sequences, their convolution

is the sequence $c_0, c_1, c_2, \dots, c_{m+n}$

where

$$c_k = \sum_{i+j=k} a_i b_j$$

e.g.

$$c_3 = a_0 b_3 + a_1 b_2 + a_2 b_1 + a_3 b_0.$$

Often (a_i) is a "signal" / "image"

(b_j) is a "mask" / "Weights"

If the sequences $(a_i), (b_j)$ are encoded as polynomials

$$A(x) = a_0 + a_1 x + a_2 x^2 + \dots + a_n x^n$$

$$B(x) = b_0 + b_1 x + \dots + b_m x^m$$

then their product $A \cdot B$ is a polynomial whose coeffs are the convolution of (a_i) and (b_j) .

3. Multiplying polynomials encodes summing independent random variables.

Ex. If one rolls two standard dice what is the probability of every possible sum?

Outcome: 1 2 3 4 5 6
 probability: $\frac{1}{6} \quad \frac{1}{6} \quad \frac{1}{6} \quad \frac{1}{6} \quad \frac{1}{6} \quad \frac{1}{6}$

$$\frac{1}{6}(x + x^2 + x^3 + \dots + x^6)$$

↑ "Prob. generating function"

$$\sum_i \Pr(i) x^i$$

JF two indep rand vars. have

prob. gen. funk's A and B

$$A(x) = \sum a_i x^i = \sum \Pr(\text{"Variable A" } = i) \cdot x^i$$

$$B(x) \quad - \quad - \quad -$$

then $A(x) \cdot B(x)$ is the PGF of their sum.

$$\text{Ex. } \frac{1}{36}(x + x^2 + x^3 + x^4 + x^5 + x^6)^2$$

$$= \frac{1}{36} x^2 + \frac{2}{36} x^3 + \frac{3}{36} x^4 + \dots + \frac{2}{36} x^{11} + \frac{1}{36} x^{12}$$

