

Operating Systems

Where's the puck going?





Announcements

- Final: 12/11 @ 9AM, Barton Hall, 100 west
- Review session: 12/08, 2PM, zoom (post on Ed discussions)
- Lost sessions: thanks for using; makes me happy about my experiments
- Please fill out the course evaluations
 - Easy way to get 5%
 - Please be constructive (evaluations are for many eyes, not just me)

Taking 25 steps back!

What is an operating system, and what does it do?

A **software layer** designed with three goals:

- Enable applications to conveniently access hardware
- Manage all hardware resources
- Implement common services for applications

What does an OS do?

- Enables **convenient "abstractions"** for applications to access hardware
 - CPU: threads
 - Memory: virtual memory
 - Storage devices: files
 - Network: sockets
 - Server: collection of resources needed by an application
 - processes, VM, containers

What does an OS do?

- Enables **convenient "abstractions"** for applications to access hardware
- Manages hardware resources
 - Resource **allocation** to individual applications
 - Resource **sharing** across concurrently running applications
 - Resource isolation across concurrently running applications

What does an OS do?

- Enables convenient "abstractions" for applications to access hardware
- Manages hardware resources
- Implements common services for applications
 - Security, protection and authentication
 - Reliability
 - Communication
 - Input/output operations
 - Program execution

•

Four Fundamental OS Concepts

• Thread: Execution Context

- A single, sequential execution context
- Address space (with translation)
 - Program's view of memory is distinct from physical memory
- Process: an instance of a running program
 - Address Space + One or more Threads + ...

Protection/Isolation

- Only the "system" can access certain resources
- Combined with translation, isolates programs from each other

Threads

- A single, sequential execution context
- A virtual core: provides illusion of infinite cores
 - Enables efficient multiplexing of physical cores...
 - ...across concurrently running applications
- Challenges in designing virtual cores
 - Scheduling, synchronization
- The OS provides protection/isolation at process granularity
 - Each thread has its own state
 - Can access other threads' state (within the same process)

CPU scheduling

- Many different possible scheduling mechanisms
- FIFO, SJF, EDF, RR, SRTF
 - Some are preemptive, some are not preemptive
 - No one-size-fits-all solution
- Our focus: understanding tradeoffs (pros and cons) of each mechanism
 - And using the insights to build a near-ideal CPU scheduler
 - Very close to the Linux CFS scheduler
- Some conceptual takeaways that we studied
 - Priority scheduling can "emulate" most scheduling mechanisms
 - Priorities should be used to define physical core share
 - Rather than strictly preferential job scheduling

Synchronization

- Coordination between multiple...
 - ...threads within the same protection/isolation domain
 - ...processes and threads operating on shared data
- A hard problem
 - No "algorithm" to design a correct-by-design program
- Our focus:
 - Understanding the core challenges in synchronization
 - A suite of techniques that can be used
 - Locks, semaphores, condition variables, monitors
 - Hardware support for synchronization

Memory management

- Virtual address space: virtualizing physical memory address space
 - Enables efficient multiplexing of memory...
 - ...across concurrently running applications
- We focused on three aspects in memory management
- Efficient sharing of physical resources
 - Paging, and page replacement
- Space and time efficient address translation
 - Space efficiency: multi-level page tables
 - Time efficiency: TLB, small #levels in multi-level page tables
- Protection
 - Apps use virtual address, kernel handles physical addresses

Memory management

- Virtual memory: provides illusion of infinite memory
 - By swapping/paging data to secondary storage
 - Each program gets the illusion of having dedicated, infinite, memory
- Paging
 - Page faults
 - Page replacement mechanisms:
 - Optimal (Belady's algorithm)
 - LRU
 - Approximating LRU: The clock algorithm
 - Working set page replacement
 - Local and global page replacement

Beyond threads, processes and memory

- The OS must handle all IO devices
 - Storage devices: HDD, SSD
 - Network devices: NIC
 - Peripheral devices: mouse, keyboards, ...
 - And all buses: memory bus, I/O bus, peripheral bus
- Mechanisms: Interrupt-driven I/O, DMA
- **Devices**: Mostly SSD, brief discussion on HDD

Beyond threads, processes and memory

- The OS must handle all IO devices
 - Storage devices: HDD, SSD
 - Network devices: NIC
 - Peripheral devices: mouse, keyboards, ...
 - And all buses: memory bus, I/O bus, peripheral bus
- OS support for handling <u>storage</u> devices
 - File systems
 - contiguous, linked list, tree-based multi-level index file storage
 - consistent updates
 - Block layer
 - Device drivers (minimal discussion)

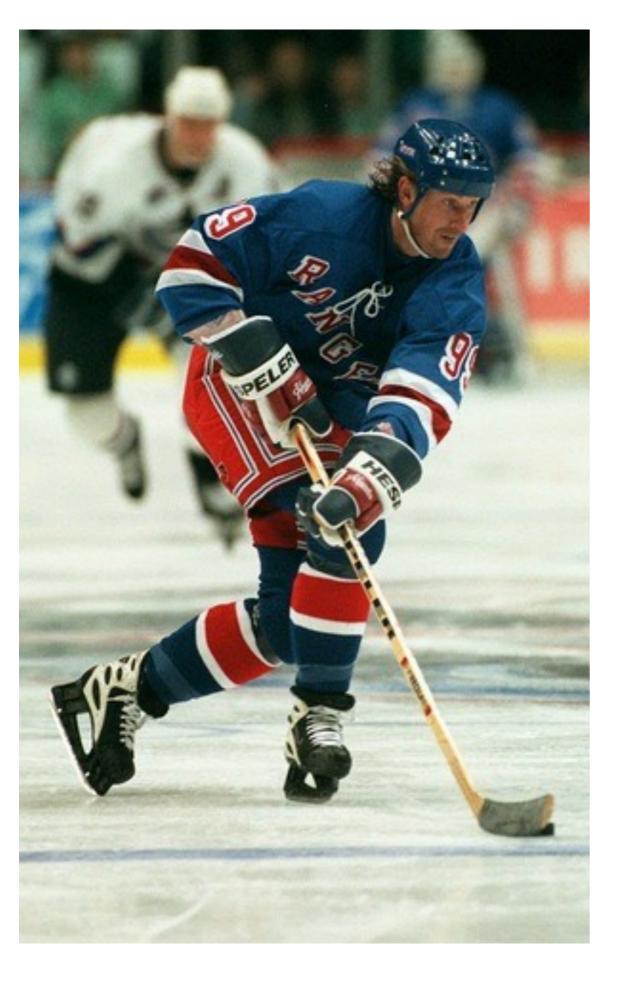
Beyond threads, processes and memory

- The OS must handle all IO devices
 - Storage devices: HDD, SSD
 - Network devices: NIC
 - Peripheral devices: mouse, keyboards, ...
 - And all buses: memory bus, I/O bus, peripheral bus

OS support for handling <u>network</u> devices

- The entire "network stack"
- End-to-end story
- Various functionalities that interact with other layers
 - Sockets and ports
 - Packet steering, and tradeoffs
 - Packet aggregation, and tradeoffs

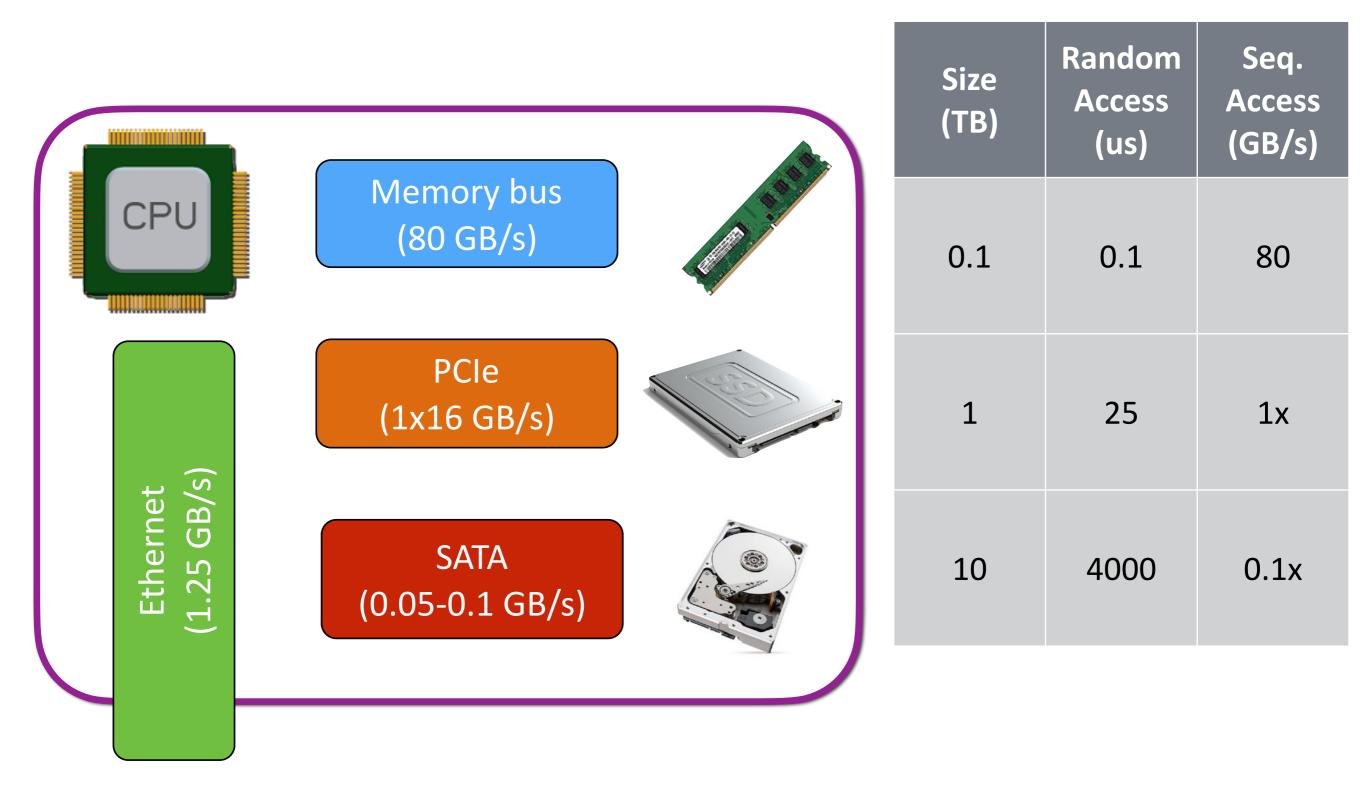
Taking 1 step forward!



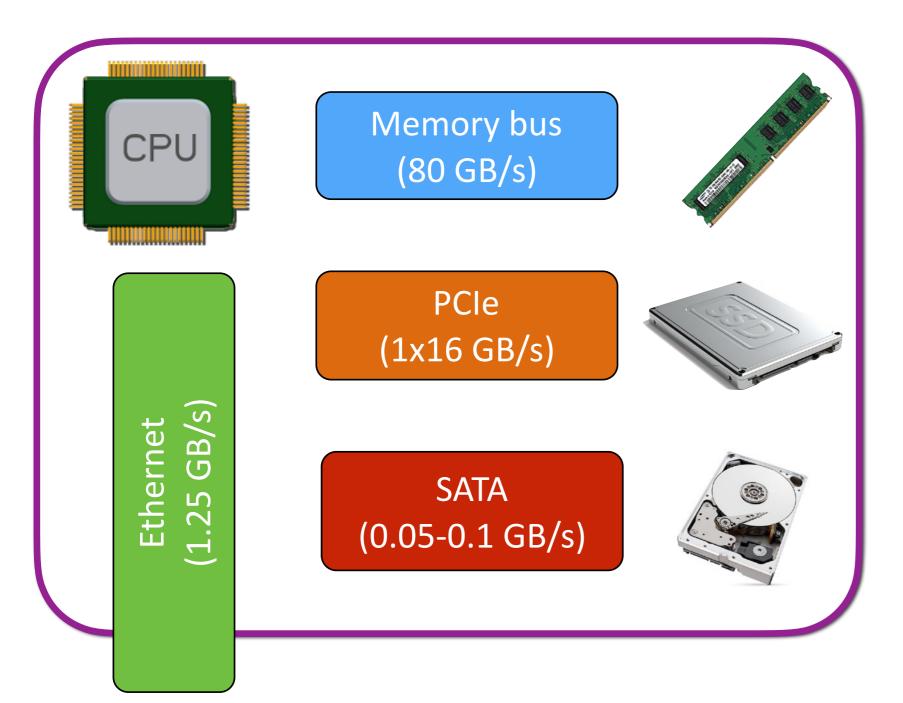
Skate where the puck's going, not where it's been!

- Walter Gretzky

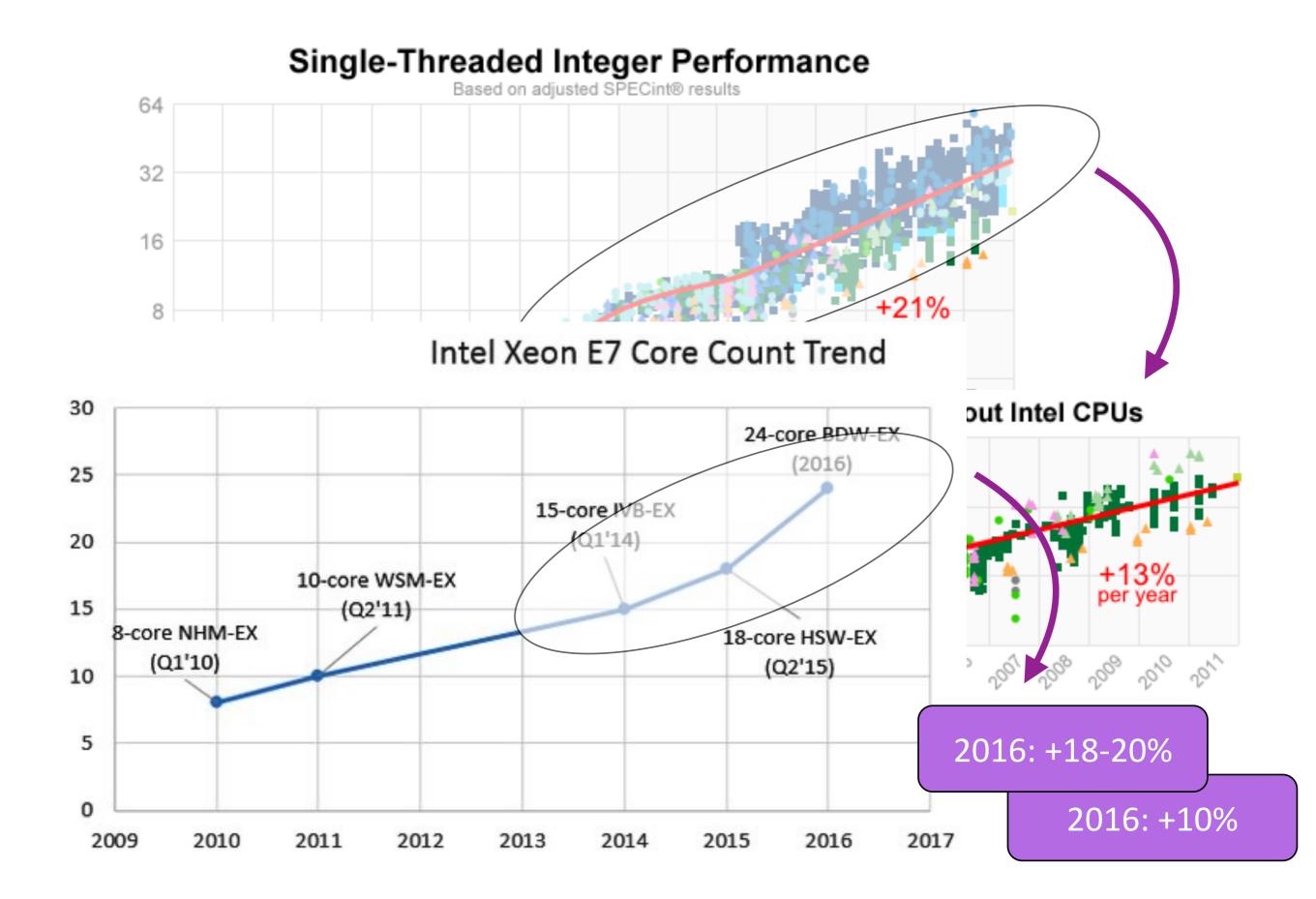
Where is the puck right now?



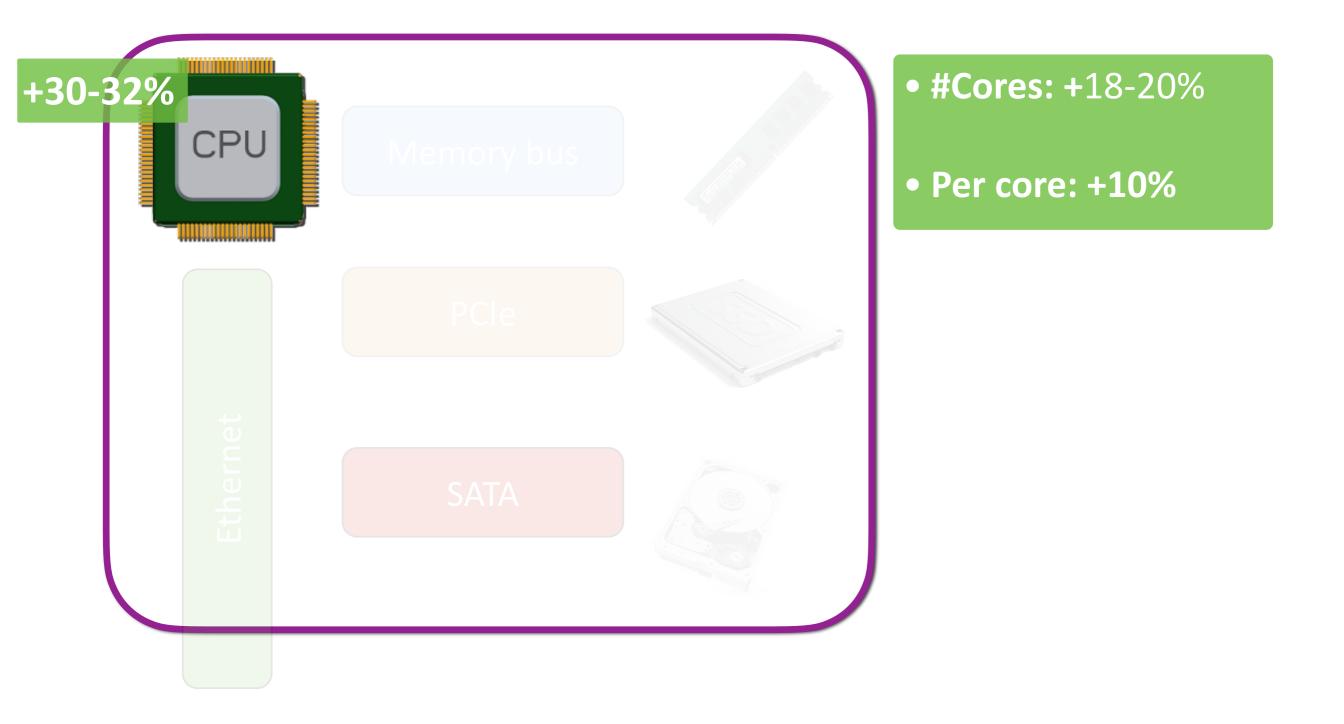
Where is the puck going?



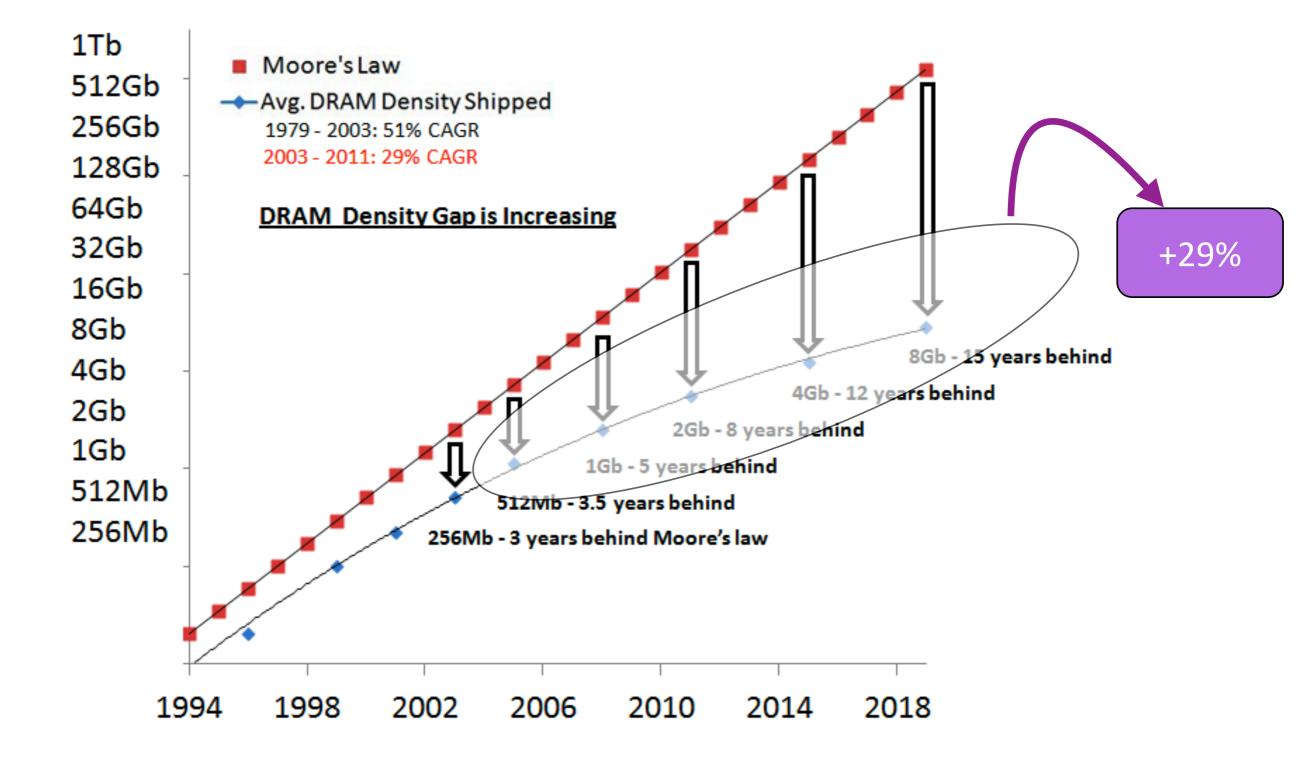
Where is the puck going? (CPU performance)



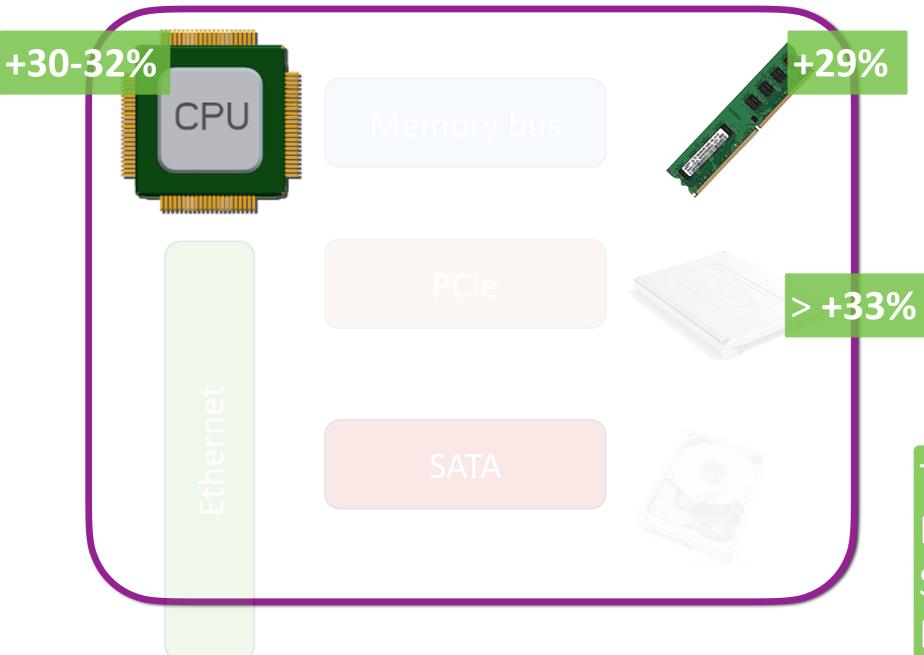
Where is the puck going?



Where is the puck going? (DRAM capacity)



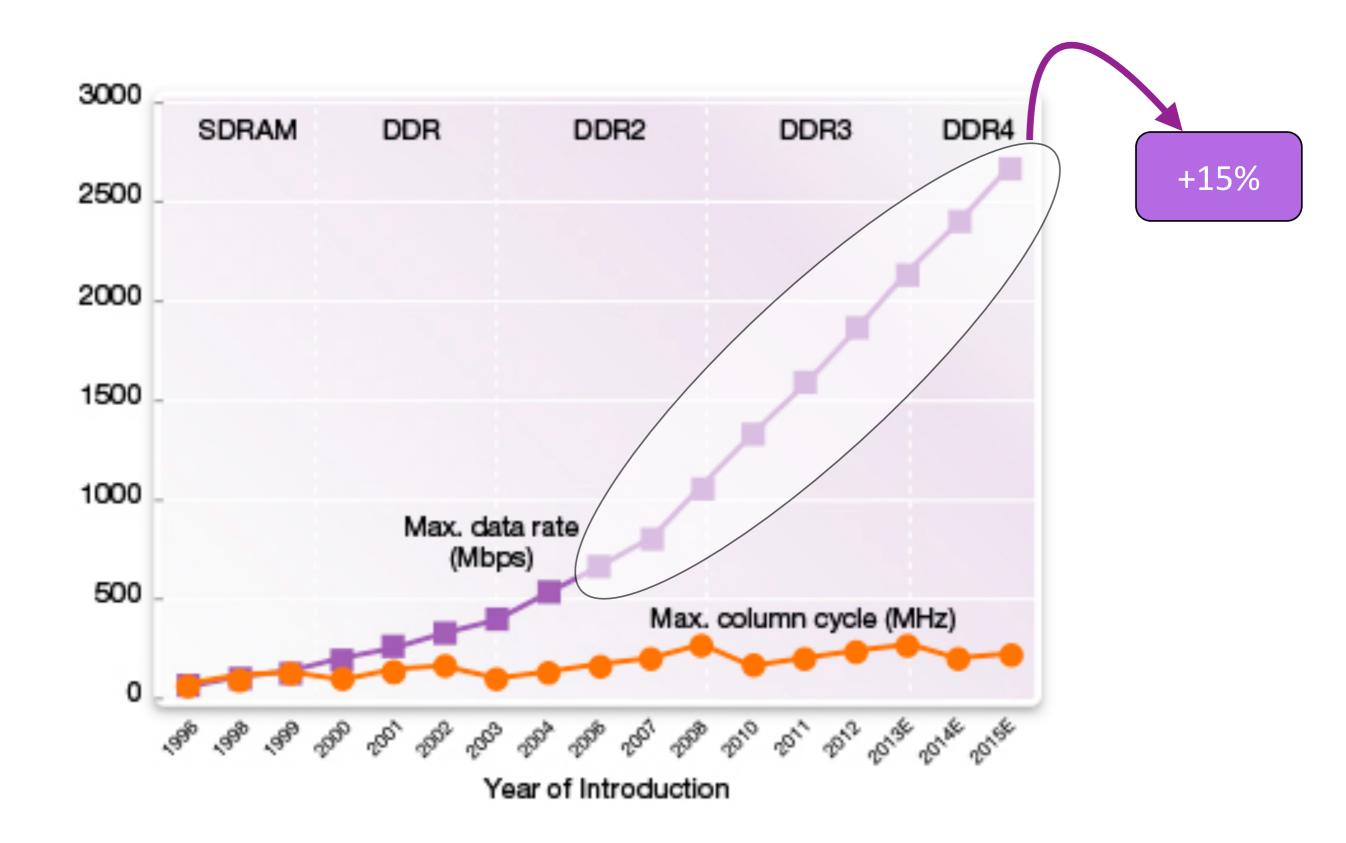
Where is the puck going?



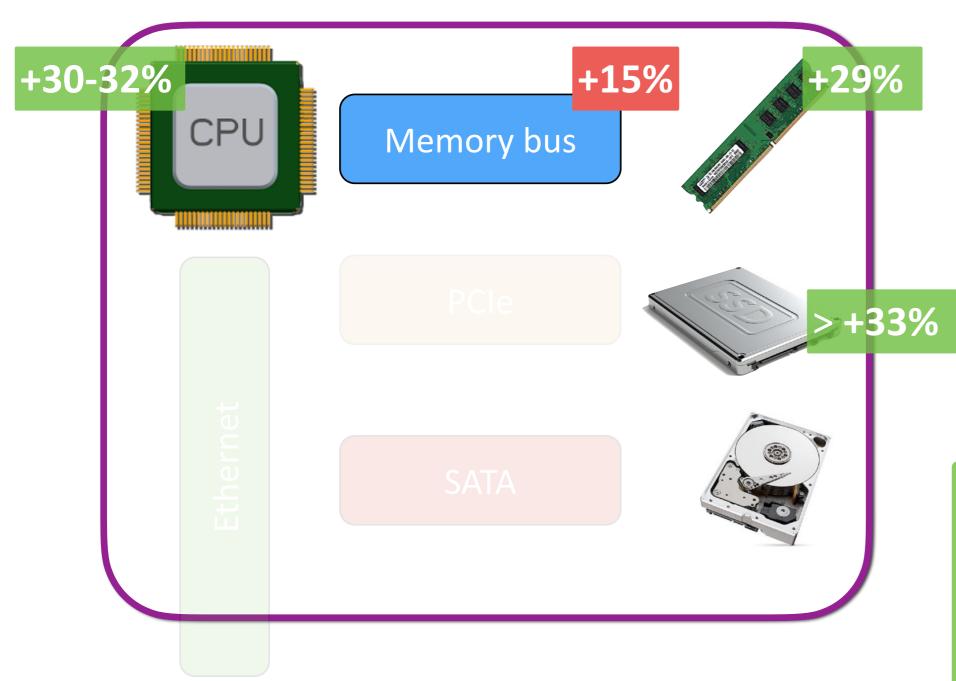
Tape is dead, Disk is tape, SSD is disk, RAM is the king!

Jim Gray

Where is the puck going? (Memory bus)



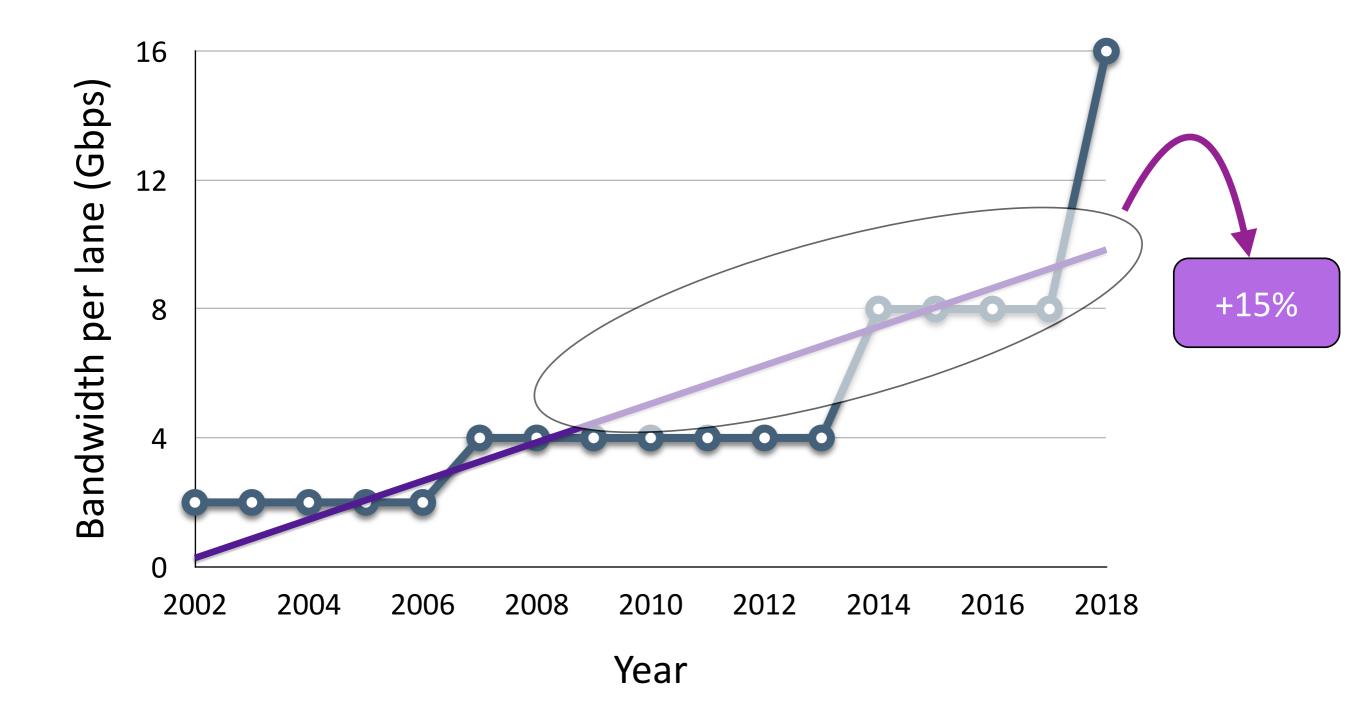
Where is the puck going?



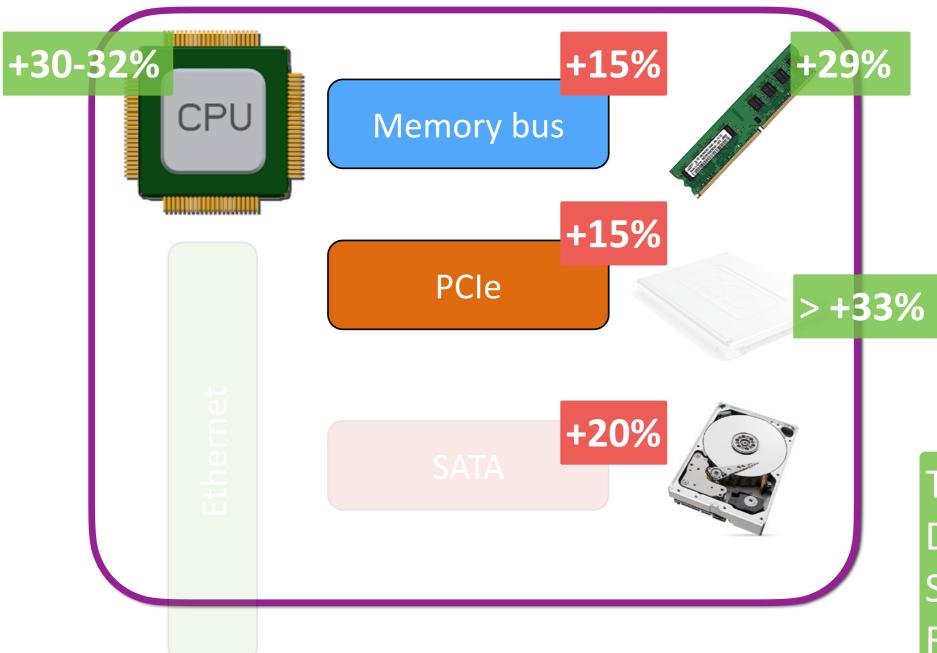
Tape is dead, Disk is tape, SSD is disk, RAM is the king!

- Jim Gray

Where is the puck going? (PCIe)

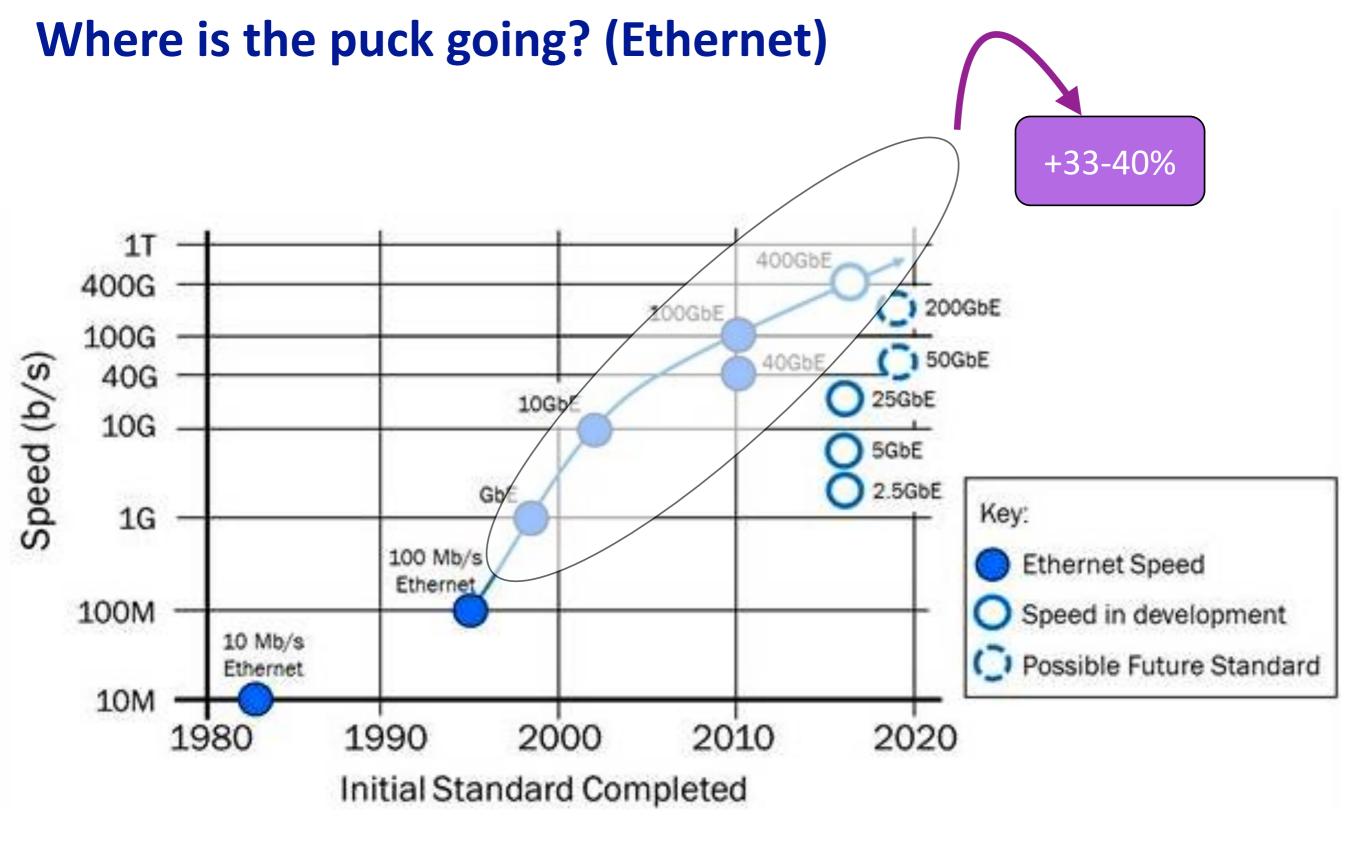


Where is the puck going?

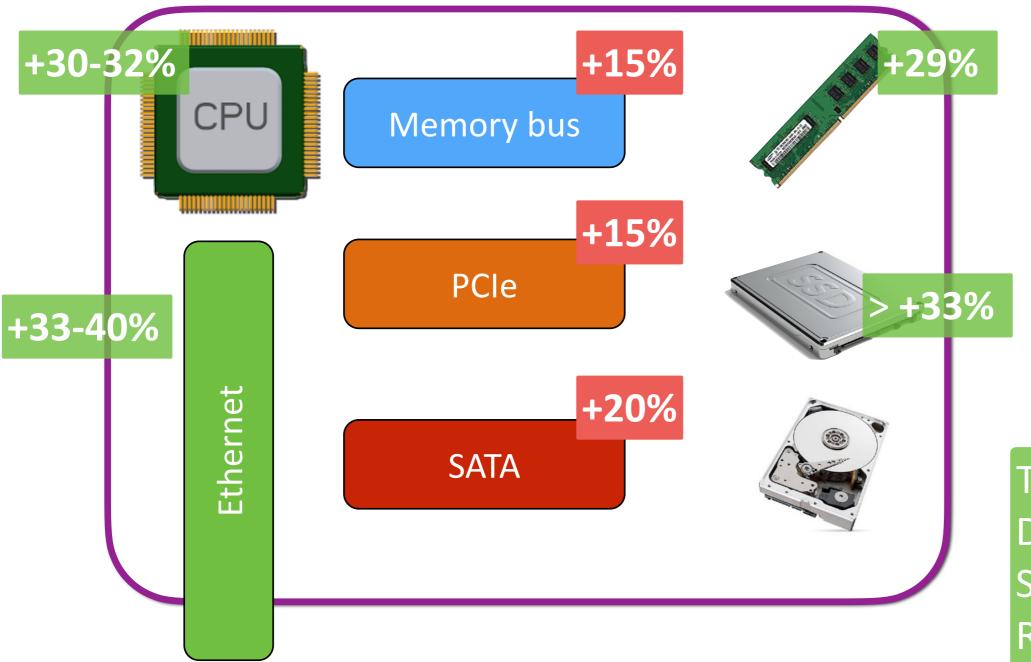


Tape is dead, Disk is tape, SSD is disk, RAM is the king!

- Jim Gray



Where is the puck going?



Tape is dead, Disk is tape, SSD is disk, RAM is the king!

- Jim Gray

Many powerful implications

- CPU is becoming the core bottleneck
 - Storage devices can achieve 10-100x higher throughput
 - NIC can transmit/receive 10-100x more packets
 - PCle can transmit/receive 10-100x more data
 - But CPU capacity is mostly stagnant
- New devices are emerging
 - New hardware accelerators: for apps that require more compute
 - FPGAs, TPUs, SmartNICs
 - Non-volatile memory devices
 - byte addressable, but persistent
 - 10x slower than main memory, 10x faster than SSD
 - RDMA NICs
 - Can read/write to memory on other servers without CPU

CPU is becoming the bottleneck

Today, CPU involved in all steps

- Running applications
- Kernel processing
- 1/0

Many of these are heavy-weight operations

- Thread and process state management
- Context switches
- Swapping and paging
- Storage access
- Network access
- Need to rethink design/optimization of each of these layers

Emergence of new devices [Compute]

- New hardware accelerators: for apps that require more compute
 - How should the OS enable sharing of accelerators?
 - How should the OS orchestrate traditional CPU and accelerator resources?
 - How should CPUs and accelerators share memory?

• Requires rethinking the abstractions developed over decades

- Threads
- Processes
- Virtual address space, and virtual memory
- Sockets

Emergence of new devices [Storage]

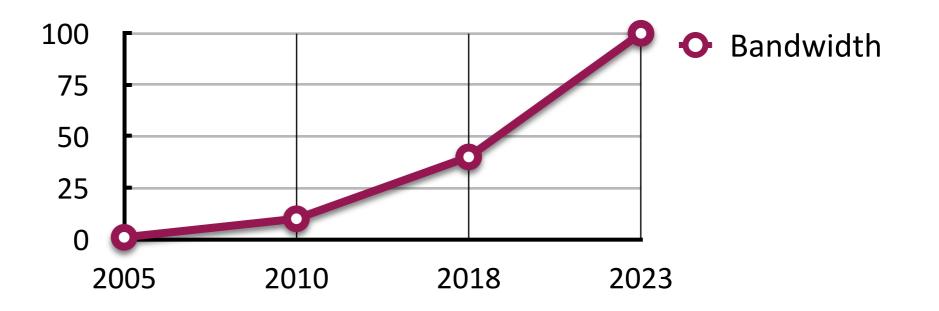
Non-volatile memory devices

- Byte-addressable—like main memory
- Persistent—like SSDs
- 10x slower than main memory, 10x faster than SSDs

Requires rethinking the abstractions developed over decades

- Virtual address space
- Virtual memory
- Page replacement

Remote Memory Faster than Local Storage



• Under zero queueing:

- Remote memory access takes less than 6.3us
- Local SSD access latency today is 25us (hardware, ignoring stack)
- Remote Direct Memory Access (RDMA) becomes feasible

Emergence of new devices [Network]

Remote Direct Memory Access

- Enables accessing remote server memory....
 - ...without involving remote server CPU
- "Kernel-bypass": CPUs can read/write data to NIC without kernel

Requires rethinking the abstractions developed over decades

- Sockets
- Protection/isolation
- Virtual address space, and virtual memory

Operating Systems are the bottleneck again!

- Lot of research in "user space designs" and kernel-bypass
 - Minimize kernel involvement
 - Low-overhead CPU scheduling
 - Lots of interesting challenges
- Lot of research in low-overhead storage stack design
 - Revisiting File systems, virtual memory, block layer, ...
 - To minimize CPU utilization, to achieve low latency and high throughput
 - Extremely interesting challenges
- Lot of research in low-overhead network stack design
 - Revisiting the many layers within the network stack
 - To minimize CPU utilization, to achieve low latency and high throughput
 - Requires rethinking host architecture, and host network
 - One of the biggest challenges faced by the OS community

Closing Thoughts

- These are exciting times for operating systems
 - The first ever since the invention of SSDs!
 - You are witness to the transformation!!!!
- And, I am glad I got the chance to introduce you to this world :-)
 - You have made me a better teacher!!!!
 - Thank you.
- Wherever you end up:
 - Please remember me
 - Say hello if you see me
 - Remember, there is nothing more important than
 - Knowing the fundamentals!!!!
 - Being happy!!!!