


Relational Query Optimization

Chapter 15


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Overview of Query Optimization

- ❖ Plan: Tree of R.A. ops, with choice of alg for each op.
 - Each operator typically implemented using a 'pull' interface: when an operator is 'pulled' for the next output tuples, it 'pulls' on its inputs and computes them.
- ❖ Two main issues:
 - For a given query, what plans are considered?
 - ◆ Algorithm to search plan space for cheapest (estimated) plan.
 - How is the cost of a plan estimated?
- ❖ Ideally: Want to find best plan. Practically: Avoid worst plans!
- ❖ We will study the System R approach.


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Highlights of System R Optimizer

- ❖ Impact:
 - Most widely used currently; works well for < 10 joins.
- ❖ Cost estimation: Approximate art at best.
 - Statistics, maintained in system catalogs, used to estimate cost of operations and result sizes.
 - Considers combination of CPU and I/O costs.
- ❖ Plan Space: Too large, must be pruned.
 - Only the space of *left-deep plans* is considered.
 - ◆ Left-deep plans allow output of each operator to be *pipelined* into the next operator without storing it in a temporary relation.
 - Cartesian products avoided.

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Schema for Examples

Sailors (*sid*: integer, *sname*: string, *rating*: integer, *age*: real)
 Reserves (*sid*: integer, *bid*: integer, *day*: dates, *rname*: string)

- ❖ Similar to old schema; *rname* added for variations.
- ❖ Reserves:
 - Each tuple is 40 bytes long, 100 tuples per page, 1000 pages.
- ❖ Sailors:
 - Each tuple is 50 bytes long, 80 tuples per page, 500 pages.

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Motivating Example

SELECT S.sname
 FROM Reserves R, Sailors S
 WHERE R.sid=S.sid AND
 R.bid=100 AND S.rating>5

❖ Cost: 500+500*1000 I/Os

❖ By no means the worst plan!

❖ Misses several opportunities: selections could have been 'pushed' earlier, no use is made of any available indexes, etc.

❖ Goal of optimization: To find more efficient plans that compute the same answer.

RA Tree: Π_{sname}

$\sigma_{bid=100} \wedge \sigma_{rating > 5}$

$\bowtie_{sid=sid}$

Reserves Sailors


Plan: Π_{sname} (On-the-fly)

$\sigma_{bid=100} \wedge \sigma_{rating > 5}$ (On-the-fly)

$\bowtie_{sid=sid}$ (Simple Nested Loops)

Reserves Sailors

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Alternative Plans 1 (No Indexes)

Π_{sname} (On-the-fly)

$\bowtie_{sid=sid}$ (Sort-Merge Join)

$\sigma_{bid=100}$ (Scan; write to temp T1)

$\sigma_{rating > 5}$ (Scan; write to temp T2)

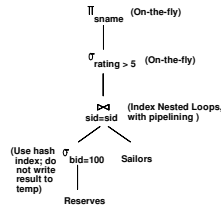
Reserves Sailors

- ❖ Main difference: push selects.
- ❖ With 5 buffers, cost of plan:
 - Scan Reserves (1000) + write temp T1 (10 pages, if we have 100 boats, uniform distribution).
 - Scan Sailors (500) + write temp T2 (250 pages, if we have 10 ratings).
 - Sort T1 (2*2*10), sort T2 (2*3*250), merge (10+250)
 - Total: 3560 page I/Os.
- ❖ If we used BNL join, join cost = 10+4*250, total cost = 2770.
- ❖ If we 'push' projections, T1 has only *sid*, T2 only *sid* and *sname*:
 - T1 fits in 3 pages, cost of BNL drops to under 250 pages, total < 2000.

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Alternative Plans 2 With Indexes

- With clustered index on *bid* of Reserves, we get $100,000/100 = 1000$ tuples on $1000/100 = 10$ pages.
- INL with *pipelining* (outer is not materialized).
 - Projecting out unnecessary fields from outer doesn't help.
- Join column *sid* is a key for Sailors.
 - At most one matching tuple, unclustered index on *sid* OK.
- Decision not to push *rating>5* before the join is based on availability of *sid* index on Sailors.
- Cost: Selection of Reserves tuples (10 I/Os); for each, must get matching Sailors tuple (1000*1.2); total 1210 I/Os.



Query Blocks: Units of Optimization

- An SQL query is parsed into a collection of *query blocks*, and these are optimized one block at a time.
- Nested blocks are usually treated as calls to a subroutine, made once per outer tuple. (This is an over-simplification, but serves for now.)
- For each block, the plans considered are:
 - All available access methods, for each reln in FROM clause.
 - All *left-deep join trees* (i.e., all ways to join the relations one-at-a-time, with the inner reln in the FROM clause, considering all reln permutations and join methods.)

```
SELECT S.sname
FROM Sailors S
WHERE S.age IN
  (SELECT MAX (S2.age)
   FROM Sailors S2
   GROUP BY S2.rating)
```

Outer block Nested block

Cost Estimation

- For each plan considered, must estimate cost:
 - Must estimate *cost* of each operation in plan tree.
 - Depends on input cardinalities.
 - We've already discussed how to estimate the cost of operations (sequential scan, index scan, joins, etc.)
 - Must estimate *size of result* for each operation in tree!
 - Use information about the input relations.
 - For selections and joins, assume independence of predicates.
- We'll discuss the System R cost estimation approach.
 - Very inexact, but works ok in practice.
 - More sophisticated techniques known now.

Statistics and Catalogs

- Need information about the relations and indexes involved. *Catalogs* typically contain at least:
 - # tuples (NTuples) and # pages (NPages) for each relation.
 - # distinct key values (NKeys) and NPages for each index.
 - Index height, low/high key values (Low/High) for each tree index.
- Catalogs updated periodically.
 - Updating whenever data changes is too expensive; lots of approximation anyway, so slight inconsistency ok.
- More detailed information (e.g., histograms of the values in some field) are sometimes stored.

Size Estimation and Reduction Factors

- Consider a query block:

```
SELECT attribute list
FROM relation list
WHERE term1 AND ... AND termk
```
- Maximum # tuples in result is the product of the cardinalities of relations in the FROM clause.
- Reduction factor (RF)* associated with each *term* reflects the impact of the *term* in reducing result size. *Result cardinality* = Max # tuples * product of all RF's.
 - Implicit assumption that *terms* are independent!
 - Term *col=value* has RF $1/NKeys(I)$, given index I on *col*
 - Term *col1=col2* has RF $1/MAX(NKeys(I1), NKeys(I2))$
 - Term *col>value* has RF $(High(I)-value)/(High(I)-Low(I))$

Relational Algebra Equivalences

- Allow us to choose different join orders and to 'push' selections and projections ahead of joins.
- Selections: $\sigma_{c_1 \wedge \dots \wedge c_n}(R) \equiv \sigma_{c_1}(\dots \sigma_{c_n}(R))$ (Cascade)
- $\sigma_{c_1}(\sigma_{c_2}(R)) \equiv \sigma_{c_2}(\sigma_{c_1}(R))$ (Commute)
- Projections: $\pi_{a_1}(R) \equiv \pi_{a_1}(\dots (\pi_{a_n}(R)))$ (Cascade)
- Joins: $R \bowtie (S \bowtie T) \equiv (R \bowtie S) \bowtie T$ (Associative)
- $(R \bowtie S) \equiv (S \bowtie R)$ (Commute)
- ⚡ Show that: $R \bowtie (S \bowtie T) \equiv (T \bowtie R) \bowtie S$

More Equivalences

- ❖ A projection commutes with a selection that only uses attributes retained by the projection.
- ❖ Selection between attributes of the two arguments of a cross-product converts cross-product to a join.
- ❖ A selection on just attributes of R commutes with $R \bowtie S$. (i.e., $\sigma(R \bowtie S) \equiv \sigma(R) \bowtie S$)
- ❖ Similarly, if a projection follows a join $R \bowtie S$, we can 'push' it by retaining only attributes of R (and S) that are needed for the join or are kept by the projection.

Enumeration of Alternative Plans

- ❖ There are two main cases:
 - Single-relation plans
 - Multiple-relation plans
- ❖ For queries over a single relation, queries consist of a combination of selects, projects, and aggregate ops:
 - Each available access path (file scan / index) is considered, and the one with the least estimated cost is chosen.
 - The different operations are essentially carried out together (e.g., if an index is used for a selection, projection is done for each retrieved tuple, and the resulting tuples are *pipelined* into the aggregate computation).

Cost Estimates for Single-Relation Plans

- ❖ Index I on primary key matches selection:
 - Cost is $Height(I)+1$ for a B+ tree, about 1.2 for hash index.
 - ❖ Clustered index I matching one or more selects:
 - $(NPages(I)+NPages(R)) * \text{product of RF's of matching selects.}$
 - ❖ Non-clustered index I matching one or more selects:
 - $(NPages(I)+NTuples(R)) * \text{product of RF's of matching selects.}$
 - ❖ Sequential scan of file:
 - $NPages(R)$.
- ⚠ **Note:** Typically, no duplicate elimination on projections! (Exception: Done on answers if user says DISTINCT.)

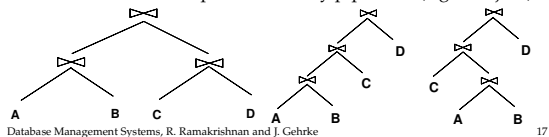
Example

```
SELECT S.sid
FROM Sailors S
WHERE S.rating=8
```

- ❖ If we have an index on *rating*:
 - $(1/NKeys(I)) * NTuples(R) = (1/10) * 40000$ tuples retrieved.
 - Clustered index: $(1/NKeys(I)) * (NPages(I)+NPages(R)) = (1/10) * (50+500)$ pages are retrieved. (This is the *cost*.)
 - Unclustered index: $(1/NKeys(I)) * (NPages(I)+NTuples(R)) = (1/10) * (50+40000)$ pages are retrieved.
- ❖ If we have an index on *sid*:
 - Would have to retrieve all tuples/pages. With a clustered index, the cost is $50+500$, with unclustered index, $50+40000$.
- ❖ Doing a file scan:
 - We retrieve all file pages (500).

Queries Over Multiple Relations

- ❖ Fundamental decision in System R: only left-deep join trees are considered.
 - As the number of joins increases, the number of alternative plans grows rapidly; we need to restrict the search space.
 - Left-deep trees allow us to generate all *fully pipelined* plans.
 - ♦ Intermediate results not written to temporary files.
 - ♦ Not all left-deep trees are fully pipelined (e.g., SM join).



Enumeration of Left-Deep Plans

- ❖ Left-deep plans differ only in the order of relations, the access method for each relation, and the join method for each join.
- ❖ Enumerated using N passes (if N relations joined):
 - Pass 1: Find best 1-relation plan for each relation.
 - Pass 2: Find best way to join result of each 1-relation plan (as outer) to another relation. (All 2-relation plans.)
 - Pass N: Find best way to join result of a (N-1)-relation plan (as outer) to the N'th relation. (All N-relation plans.)
- ❖ For each subset of relations, retain only:
 - Cheapest plan overall, plus
 - Cheapest plan for each *interesting order* of the tuples.

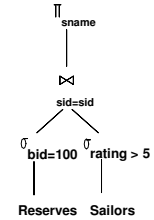
Enumeration of Plans (Contd.)

- ❖ ORDER BY, GROUP BY, aggregates etc. handled as a final step, using either an 'interestingly ordered' plan or an additional sorting operator.
- ❖ An N-1 way plan is not combined with an additional relation unless there is a join condition between them, unless all predicates in WHERE have been used up.
 - i.e., avoid Cartesian products if possible.
- ❖ In spite of pruning plan space, this approach is still exponential in the # of tables.

Example

Sailors:
B+ tree on *rating*
Hash on *sid*

Reserves:
B+ tree on *bid*



Pass 1:

- *Sailors*: B+ tree matches *rating > 5*, and is probably cheapest. However, if this selection is expected to retrieve a lot of tuples, and index is unclustered, file scan may be cheaper.
 - ♦ Still, B+ tree plan kept (because tuples are in *rating* order).
- *Reserves*: B+ tree on *bid* matches *bid=500*; cheapest.

Pass 2:

- We consider each plan retained from Pass 1 as the outer, and consider how to join it with the (only) other relation.
 - ♦ e.g., *Reserves as outer*: Hash index can be used to get *Sailors* tuples that satisfy *sid = outer tuple's sid* value.

Nested Queries

- ❖ Nested block is optimized independently, with the outer tuple considered as providing a selection condition.
- ❖ Outer block is optimized with the cost of 'calling' nested block computation taken into account.
- ❖ Implicit ordering of these blocks means that some good strategies are not considered. *The non-nested version of the query is typically optimized better.*

```
SELECT S.sname
FROM Sailors S
WHERE EXISTS
  (SELECT *
   FROM Reserves R
   WHERE R.bid=103
   AND R.sid=S.sid)
```

Nested block to optimize:

```
SELECT *
FROM Reserves R
WHERE R.bid=103
AND S.sid= outer value
```

Equivalent non-nested query:

```
SELECT S.sname
FROM Sailors S, Reserves R
WHERE S.sid=R.sid
AND R.bid=103
```

Summary

- ❖ Query optimization is an important task in a relational DBMS.
- ❖ Must understand optimization in order to understand the performance impact of a given database design (relations, indexes) on a workload (set of queries).
- ❖ Two parts to optimizing a query:
 - Consider a set of alternative plans.
 - ♦ Must prune search space; typically, left-deep plans only.
 - Must estimate cost of each plan that is considered.
 - ♦ Must estimate size of result and cost for each plan node.
 - ♦ *Key issues*: Statistics, indexes, operator implementations.

Summary (Contd.)

- ❖ Single-relation queries:
 - All access paths considered, cheapest is chosen.
 - *Issues*: Selections that *match* index, whether index key has all needed fields and/or provides tuples in a desired order.
- ❖ Multiple-relation queries:
 - All single-relation plans are first enumerated.
 - ♦ Selections/projections considered as early as possible.
 - Next, for each 1-relation plan, all ways of joining another relation (as inner) are considered.
 - Next, for each 2-relation plan that is 'retained', all ways of joining another relation (as inner) are considered, etc.
 - At each level, for each subset of relations, only best plan for each interesting order of tuples is 'retained'.