# $1/O$

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# Big Picture: Input/Output (I/O)

How does a processor interact with its environment?

Computer System = Memory + Datapath + Control + Input + Output Keyboard Network

**Display** 

Disk

# **Putting it all together**

I/O connected with I/O Controllers high-performance interconnect: processor, memory, display lower-performance interconnect: disk, keyboard, network



## **Bus Types**

### Processor – Memory ("Front Side Bus")

- Short, fast, & wide
- Mostly fixed topology, designed as a "chipset"
	- CPU + Caches + Interconnect + Memory Controller
- I/O and Peripheral busses (PCI, SCSI, …)
	- Longer, slower, & narrower
	- Flexible topology, multiple/varied connections
	- Interoperability standards for devices
	- Connect to processor-memory bus through a bridge

# I/O Device API

Typical I/O Device API

• a set of read-only or read/write registers

### Command registers

• writing causes device to do something

#### Status registers

• reading indicates what device is doing, error codes, …

#### Data registers

- Write: transfer data to a device
- Read: transfer data from a device

Every device uses this API

### How to talk to a device?

1. Programmed I/O:

special instructions talk over special busses Specify: device, data, direction

- inb \$a, 0x64 (keyboard status register)
- outb \$a, 0x60 (keyboard data register)
- Protection: only allowed in kernel mode (expensive)

# 2. Memory-Mapped I/O:

map registers into virtual address space

- Accesses to certain addresses redirected to I/O devices
- Data goes over the memory bus (faster!)
- Protection: via bits in pagetable entries
- OS+MMU+devices configure mappings

# Memory-Mapped I/O



vs. less-favored alternative = Programmed I/O:

- Syscall instructions that communicate with I/O
- Communicate via special device registers

```
Device Drivers
                               Memory Mapped I/O
Programmed I/O
                               struct kbd {
char read_kbd()
                                 char status, pad[3];
{
                                 char data, pad[3];
do {
                               };
    sleep();
                               kbd *k = mmap(...).status = jnb(0x64);
                                                     syscall
  } while(\frac{1}{\sqrt{5}}tatus & 1));
                               char read kbd()
                               \{return inb(0x60);
                                 do {
}
syscalls
                                   sleep();
                                   status = k->status;
 Clicker Question: Which is better?
                                 } while(!(status & 1));
(A) Programmed I/O 
                                 return k->data;
 (B) Memory Mapped I/O
 (C) Both have syscalls, both are bad
```
#### I/O Data Transfer How to talk to device? • Programmed I/O or Memory-Mapped I/O How to get events? • Polling or Interrupts How to transfer lots of data?  $disk$ ->cmd = READ 4K SECTOR; Very,  $disk-\lambda data = 12;$ *Very,* while (!(disk->status & 1) { } **Expensive** for  $(i = 0..4k)$  $buff[i] = disk - > data;$

### **Data Transfer**

- 1. Programmed: Device  $\leftarrow$   $\rightarrow$  CPU  $\leftarrow$   $\rightarrow$  RAM Transfer for  $(i = 1 .. n)$ 
	- CPU issues read request
	- Device puts data on bus & CPU reads into registers
	- CPU writes data to memory



2. Direct Memory Access (DMA): Device  $\leftarrow \rightarrow$  RAM

- CPU sets up DMA request
- for  $(i = 1 ... n)$ Device puts data on bus & RAM accepts it





### Programmed I/O vs Memory Mapped I/O

### Programmed I/O

- Requires special instructions
- Can require dedicated hardware interface to devices
- Protection enforced via kernel mode access to instructions
- Virtualization can be difficult

#### Memory-Mapped I/O

- Re-uses standard load/store instructions
- Re-uses standard memory hardware interface
- Protection enforced with normal memory protection scheme
- Virtualization enabled with normal memory virtualization scheme

### **Polling vs. Interrupts**

How does program learn device is ready/done?

- 1. Polling: Periodically check I/O status register
	- Common in small, cheap, or real-time embedded systems
	- + Predictable timing, inexpensive
	- Wastes CPU cycles
- 2. Interrupts: Device sends interrupt to CPU
	- Cause register identifies the interrupting device
	- Interrupt handler examines device, decides what to do
	- + Only interrupt when device ready/done
	- Forced to save CPU context (PC, SP, registers, *etc.*)
	- Unpredictable, event arrival depends on other devices' activity

Clicker Question: Which is better? (A) Polling (B) Interrupts (C) Both equally good/bad