I/O

Anne Bracy CS 3410 Computer Science

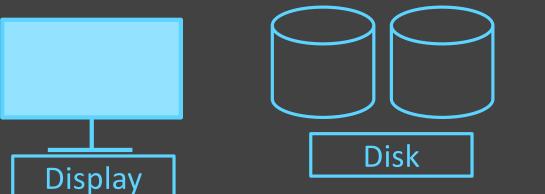
Cornell University

[K. Bala, A. Bracy, S. McKee, E. Sirer, and H. Weatherspoon]

Big Picture: Input/Output (I/O)

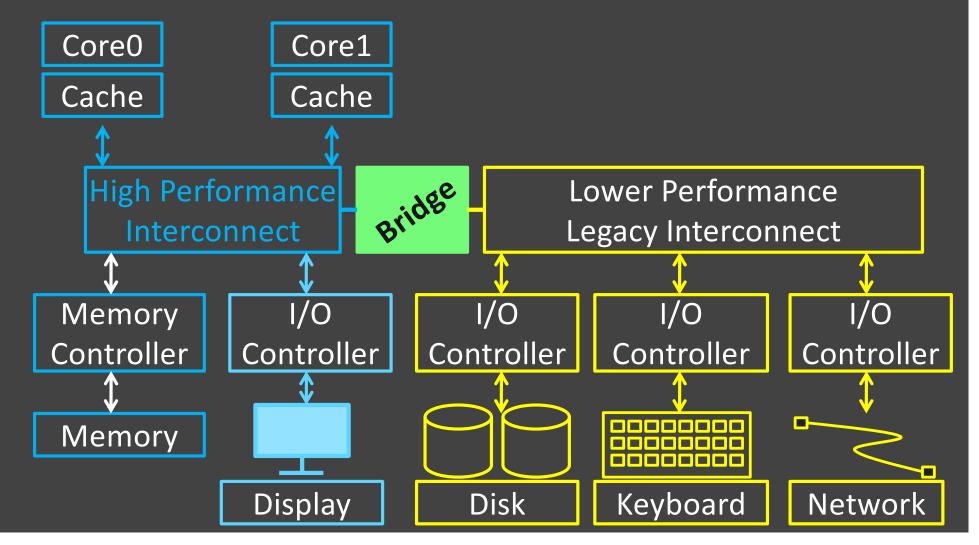
How does a processor interact with its environment?

Computer System = Memory + Datapath + Control + Input + Output Impute Input - Control + Input + Output



Putting it all together

I/O connected with I/O Controllers high-performance interconnect: processor, memory, display lower-performance interconnect: disk, keyboard, network



Bus Types

Processor - Memory ("Front Side Bus")

- Short, fast, & wide
- Mostly fixed topology, designed as a "chipset"
 - CPU + Caches + Interconnect + Memory Controller
- I/O and Peripheral busses (PCI, SCSI, ...)
 - Longer, slower, & narrower
 - Flexible topology, multiple/varied connections
 - Interoperability standards for devices
 - Connect to processor-memory bus through a bridge

I/O Device API

Typical I/O Device API

• a set of read-only or read/write registers

Command registers

writing causes device to do something

Status registers

• reading indicates what device is doing, error codes, ...

Data registers

- Write: transfer data to a device
- Read: transfer data from a device

Every device uses this API

How to talk to a device?

1. Programmed I/O:

special instructions talk over special busses Specify: device, data, direction

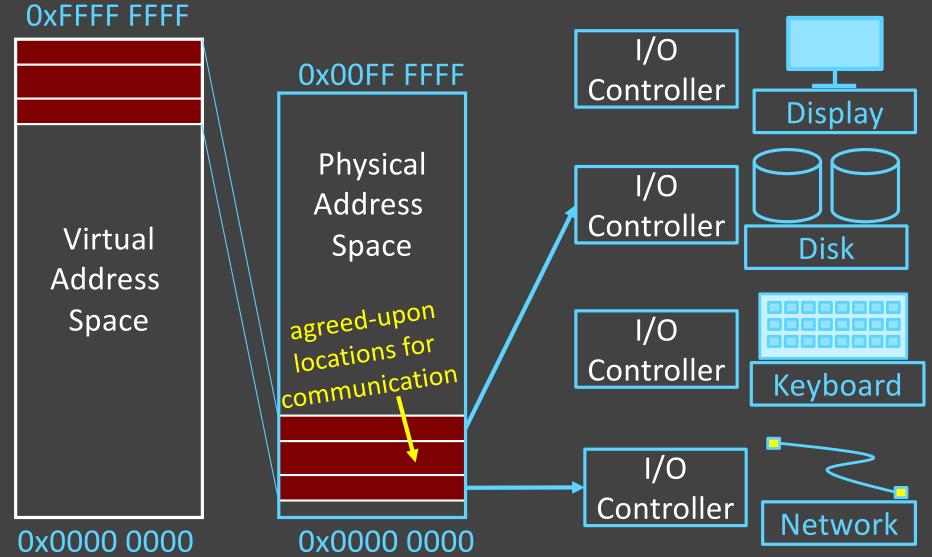
- inb \$a, 0x64 (keyboard status register)
- outb \$a, 0x60 (keyboard data register)
- Protection: only allowed in kernel mode (expensive)

2. Memory-Mapped I/O:

map registers into virtual address space

- Accesses to certain addresses redirected to I/O devices
- Data goes over the memory bus (faster!)
- Protection: via bits in pagetable entries
- OS+MMU+devices configure mappings

Memory-Mapped I/O



vs. less-favored alternative = Programmed I/O:

- Syscall instructions that communicate with I/O
- Communicate via special device registers

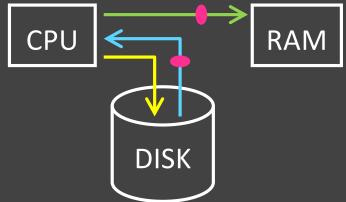
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Device Drivers
                             Memory Mapped I/O
Programmed I/O
                              struct kbd {
char read_kbd()
                                char status, pad[3];
{
                                char data, pad[3];
do {
                             };
    sleep();
                             kbd *k = mmap(...); <
    status = inb(0x64);
  } while(!(status & 1));
                             char read_kbd()
  return inb(0x60);
                                do {
} syscalls
                                  sleep();
                                  status = k->status;
Clicker Question: Which is better?
                                } while(!(status & 1));
(A) Programmed I/O
                                return k->data;
(B) Memory Mapped I/O
(C) Both have syscalls, both are bad
```

syscall

I/O Data Transfer How to talk to device? Programmed I/O or Memory-Mapped I/O How to get events? Polling or Interrupts How to transfer lots of data? disk->cmd = READ 4K SECTOR; Very, disk->data = 12; Very, while (!(disk->status & 1) { } **Expensive** for (i = 0..4k)buf[i] = disk->data;

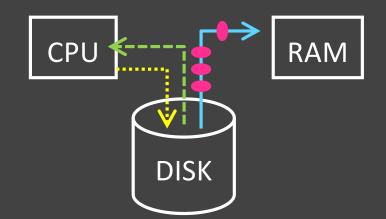
Data Transfer

- 1. Programmed: Device $\leftarrow \rightarrow$ CPU $\leftarrow \rightarrow$ RAM Transfer for (i = 1.. n)
 - CPU issues read request
 - Device puts data on bus & CPU reads into registers
 - CPU writes data to memory



2. Direct Memory Access (DMA): Device $\leftarrow \rightarrow$ RAM

- CPU sets up DMA request
- for (i = 1 ... n) Device puts data on bus & RAM accepts it
- Device interrupts CPU after done



Programmed I/O vs Memory Mapped I/O

Programmed I/O

- Requires special instructions
- Can require dedicated hardware interface to devices
- Protection enforced via kernel mode access to instructions
- Virtualization can be difficult

Memory-Mapped I/O

- Re-uses standard load/store instructions
- Re-uses standard memory hardware interface
- Protection enforced with normal memory protection scheme
- Virtualization enabled with normal memory virtualization scheme

Polling vs. Interrupts

How does program learn device is ready/done?

- 1. Polling: Periodically check I/O status register
 - Common in small, cheap, or real-time embedded systems
 - + Predictable timing, inexpensive
 - Wastes CPU cycles
- 2. Interrupts: Device sends interrupt to CPU
 - Cause register identifies the interrupting device
 - Interrupt handler examines device, decides what to do
 - + Only interrupt when device ready/done
 - Forced to save CPU context (PC, SP, registers, etc.)
 - Unpredictable, event arrival depends on other devices' activity

Clicker Question: Which is better? (A) Polling (B) Interrupts (C) Both equally good/bad