

1. Reading: K. Rosen *Discrete Mathematics and Its Applications*, 2.6
2. The main message of this lecture:

Matrices is a natural generalization of numbers. They have the same operations (addition/subtraction, multiplication/division), though with some peculiar properties: multiplication is not commutative, division is not always possible even when the denominator is not zero. Important: memorize how to multiply matrices!

Numbers here are not necessarily integers.

Definition 11.1. Matrix is a rectangular array of $m \times n$ entries. Here m is the number of rows, n - the number of columns. Notations:

$$A = \begin{pmatrix} a_{11} & a_{12} & a_{13} & \dots & a_{1n} \\ a_{21} & a_{22} & a_{23} & \dots & a_{2n} \\ a_{31} & a_{32} & a_{33} & \dots & a_{3n} \\ \dots & \dots & \dots & \dots & \dots \\ a_{m1} & a_{m2} & a_{m3} & \dots & a_{mn} \end{pmatrix} = [a_{ij}] \text{ (} i \text{ is the row index, } j \text{ is the column index)}$$

Examples:

$$\begin{pmatrix} 1 & 2 & 3 \\ 4 & 5 & 6 \end{pmatrix} \text{ (} 2 \times 3 \text{ matrix), } \begin{pmatrix} 1 & 2 \\ 3 & 4 \\ 5 & 6 \end{pmatrix} \text{ (} 3 \times 2 \text{ matrix), } \begin{pmatrix} 1 \\ 2 \\ 3 \\ 4 \end{pmatrix} \text{ (} 4 \times 1 \text{ matrix)}$$

Definition 11.2. Zero matrixes: for any given size $(m \times n)$ $\mathbf{0}$ is the matrix where all entries are equal to 0. **Addition** of matrices of the same size: add the corresponding entries $A+B = [a_{ij}+b_{ij}]$. It is clear that $A+\mathbf{0} = \mathbf{0}+A = A$. Examples: $(1\ 2\ 3\ 4)+(5\ 6\ 7\ 8) = (6\ 8\ 10\ 12)$,

$$\begin{pmatrix} 1 & 2 & 3 \\ 4 & 5 & 6 \end{pmatrix} + \begin{pmatrix} 11 & 12 & 13 \\ 14 & 15 & 16 \end{pmatrix} = \begin{pmatrix} 12 & 14 & 16 \\ 18 & 20 & 22 \end{pmatrix}, \begin{pmatrix} 1 \\ 2 \\ 3 \\ 4 \end{pmatrix} + \begin{pmatrix} -4 \\ -3 \\ -2 \\ -1 \end{pmatrix} = \begin{pmatrix} -3 \\ -1 \\ 1 \\ 3 \end{pmatrix}$$

Definition 11.3. Multiplication of matrices A ($m \times n$) and B ($n \times k$). The product of A , B is $A \cdot B = C$ ($m \times k$), where C is obtained by the *row-column rule*: to get c_{ij} multiply the corresponding entries from i th row of A and from j th column of B , and take the sum of the products.

$$c_{ij} = a_{i1}b_{1j} + a_{i2}b_{2j} + a_{i3}b_{3j} + \dots + a_{in}b_{nj} = \sum_{l=1}^n a_{il}b_{lj}$$

Examples:

$$(1\ 2\ 3\ 4) \cdot \begin{pmatrix} -4 \\ -3 \\ -2 \\ -1 \end{pmatrix} = 1 \cdot (-4) + 2 \cdot (-3) + 3 \cdot (-2) + 4 \cdot (-1) = -20$$

$$\begin{pmatrix} 1 & 2 & 3 \\ 4 & 5 & 6 \\ 0 & 1 & 0 \end{pmatrix} \cdot \begin{pmatrix} -6 & 5 \\ -4 & 3 \\ -2 & 1 \end{pmatrix} = \begin{pmatrix} 1 \cdot (-6) + 2 \cdot (-4) + 3 \cdot (-2) & 1 \cdot 5 + 2 \cdot 3 + 3 \cdot 1 \\ 4 \cdot (-6) + 5 \cdot (-4) + 6 \cdot (-2) & 4 \cdot 5 + 5 \cdot 3 + 6 \cdot 1 \\ 0 \cdot (-6) + 1 \cdot (-4) + 0 \cdot (-2) & 0 \cdot 5 + 1 \cdot 3 + 0 \cdot 1 \end{pmatrix} = \begin{pmatrix} -20 & 14 \\ -56 & 41 \\ -4 & 3 \end{pmatrix}$$

Matrix multiplication of A ($m \times n$) and B ($n \times m$) is not commutative, i.e., generally speaking, $A \cdot B \neq B \cdot A$. Indeed, even the sizes of $A \cdot B$ ($m \times m$) and $B \cdot A$ ($n \times n$) do not match. The same holds true even for square A, B . Example:

$$\begin{pmatrix} 0 & 1 \\ 0 & 0 \end{pmatrix} \cdot \begin{pmatrix} 1 & 0 \\ 0 & 0 \end{pmatrix} = \begin{pmatrix} 0 & 0 \\ 0 & 0 \end{pmatrix}, \text{ but } \begin{pmatrix} 1 & 0 \\ 0 & 0 \end{pmatrix} \cdot \begin{pmatrix} 0 & 1 \\ 0 & 0 \end{pmatrix} = \begin{pmatrix} 0 & 1 \\ 0 & 0 \end{pmatrix}$$

The former of those two examples show that $A \cdot B = \mathbf{0}$ does not yield that $A = \mathbf{0} \vee B = \mathbf{0}$.

There is no cancellation rule of matrix multiplication: $A \cdot B = A \cdot C$ and $A \neq \mathbf{0}$ does not necessarily yield $B = C$. Example:

$$\begin{pmatrix} 0 & 1 \\ 0 & 0 \end{pmatrix} \cdot \begin{pmatrix} 1 & 0 \\ 0 & 0 \end{pmatrix} = \begin{pmatrix} 0 & 1 \\ 0 & 0 \end{pmatrix} \cdot \begin{pmatrix} 0 & 0 \\ 0 & 0 \end{pmatrix} = \begin{pmatrix} 0 & 0 \\ 0 & 0 \end{pmatrix}, \text{ but } \begin{pmatrix} 1 & 0 \\ 0 & 0 \end{pmatrix} \neq \begin{pmatrix} 0 & 0 \\ 0 & 0 \end{pmatrix}$$

Complexity of the multiplication of ($m \times n$) by ($n \times k$) matrices = the number of individual additions or multiplication. In the standard algorithm for each entry $c_{ij} = a_{i1}b_{1j} + a_{i2}b_{2j} + \dots + a_{in}b_{nj}$ there are n multiplications and $n - 1$ additions. Total: mnk multiplications and $m(n-1)k$ additions, is $O(n^3)$ for ($n \times n$) matrices. There are much faster matrix multiplication algorithms than the standard one given by the definition above.

Example: Assume that we need to multiply three matrices A (10×10), B (10×100) and C (100×1). Since the matrix multiplication is associative, $A \cdot (B \cdot C) = (A \cdot B) \cdot C$, and the final result does not really depend on what multiplication is performed first. However, the complexity does! If we perform $A \cdot B$ first, we first spend $10 \cdot 10 \cdot 100$ number multiplications, and then $10 \cdot 100 \cdot 1$ number multiplications, which brings the total to $10000 + 1000 = 11000$. If $B \cdot C$ is taken first, then the total number of multiplications is $10 \cdot 100 \cdot 1 + 10 \cdot 10 \cdot 1 = 1100$ which is 1/10 of the previous estimate.

Definition 11.4. Transpose of $[a_{ij}]^T = [a_{ji}]$. Example: $\begin{pmatrix} 1 & 2 & 3 \\ 4 & 5 & 6 \end{pmatrix}^T = \begin{pmatrix} 1 & 4 \\ 2 & 5 \\ 3 & 6 \end{pmatrix}$.

A is **symmetric** if $A^T = A$ (note that such an A should be ($n \times n$)).

Definition 11.5. Identity matrix $I_n = \begin{pmatrix} 1 & 0 & \dots & 0 \\ 0 & 1 & \dots & 0 \\ \cdot & \cdot & \dots & \cdot \\ 0 & 0 & \dots & 1 \end{pmatrix}$. $A \cdot I_n = I_n \cdot A = A$ holds for any

$n \times n$ matrix A . The **power** of a square matrix A is defined by $A^0 = I_n$, $A^k = \underbrace{A \cdot A \cdot \dots \cdot A}_{k \text{ times}}$.

Definition 11.6 For 0-1 matrices $A = [a_{ij}]$ and $B = [b_{ij}]$ we can define

$$\begin{aligned} A \wedge B &= [a_{ij} \wedge b_{ij}] \text{ (meet),} & A \vee B &= [a_{ij} \vee b_{ij}] \text{ (join),} \\ A \odot B &= [(a_{i1} \wedge b_{1j}) \vee \dots \vee (a_{in} \wedge b_{nj})] \text{ (boolean product),} \\ A^{[0]} &= I_n, \quad A^{[r+1]} = A^{[r]} \odot A \text{ (boolean power).} \end{aligned}$$

Homework assignments. (due Friday 02/23).

11A:Rosen2.6-2b; 11B:Rosen2.6-4c; 11C:Rosen2.6-24a; 11D:Rosen2.6-28