

1. Reading: K. Rosen *Discrete Mathematics and Its Applications*, 2.5
2. The main message of this lecture:

**Primes are atoms of arithmetic with many striking properties. Prime factorization is very hard in practice: we can use this observation to encode messages and feel safe when an encryption key becomes public.**

Some useful math first. As before, all the numbers here are integers.

**Theorem 10.1.** *For all  $a, b > 0$  there exist  $x, y$  such that  $ax + by = \gcd(a, b)$ .*

**Proof.** By example, which is here as good as a general case. Consider  $a = 111, b = 45$ . Run the Euclidean algorithm:

$$\begin{aligned} 111 &= 2 \cdot 45 + 21 &\Rightarrow 21 &= 111 - 2 \cdot 45 \\ 45 &= 2 \cdot 21 + 3 &\Rightarrow 3 &= 45 - 2 \cdot 21 \\ 21 &= 7 \cdot 3 &\Rightarrow \gcd(111, 45) &= 3. \end{aligned}$$

Now walk these computations backward:  $3 = 45 - 2 \cdot 21 = 45 - 2 \cdot (111 - 2 \cdot 45) = 45 - 2 \cdot 111 + 4 \cdot 45 = 5 \cdot 45 - 2 \cdot 111$ . One can see easily, that  $\gcd(a, b)$  is a linear combination of each pair of remainders appearing in the process of execution of the algorithm.

**Corollary 10.2.** *If  $a, b$  are relatively prime, then there are  $x, y$ , such that  $ax + by = 1$ .*

Example:  $\gcd(101, 45) = 1$ . Find  $x, y$  such that  $101 \cdot x + 45 \cdot y = 1$ . Use the general method from 10.1. By the Euclidean Algorithm,  $101 = 2 \cdot 45 + 11, 45 = 4 \cdot 11 + 1$ . Walking backwards:  $1 = 45 - 4 \cdot 11 = 45 - 4(101 - 2 \cdot 45) = 45 - 4 \cdot 101 + 8 \cdot 45 = 45 \cdot 9 - 101 \cdot 4, x = -4, y = 9$ .

An equation  $ax \equiv b \pmod{m}$  is called **linear congruence**. Example:  $3x \equiv 2 \pmod{5}$ , solution  $x = ?$ . Let us try some  $x$ 's:  $3 \cdot 0 \equiv 0 \pmod{5}, 3 \cdot 1 \equiv 3 \pmod{5}, 3 \cdot 2 \equiv 1 \pmod{5}, 3 \cdot 3 \equiv 4 \pmod{5}, 3 \cdot 4 \equiv 2 \pmod{5}$ . Thus  $x = 4$  is a solution, as well as any number  $4 + 5k$ . Such a method is practical for small  $m$ 's: try the numbers from 0 to  $m - 1$ . Sometimes we are not lucky:  $2x \equiv 1 \pmod{4}$  has no solutions, since  $2x$  is always even, i.e. is 0 or 2 (mod 4).

**Theorem 10.3.** *If  $\gcd(a, m) = 1$  then  $ax \equiv 1 \pmod{m}$  has a solution.*

**Proof.** By 10.2, find  $x, y$  such that  $ax + my = 1$ . Then  $x$  is a solution:  $ax \equiv ax + my \equiv 1 \pmod{m}$ . Example: to solve  $45x \equiv 1 \pmod{101}$  use the above example  $1 = 45 \cdot 9 - 101 \cdot 4 \equiv 45 \cdot 9 \pmod{101}, x = 9$ .

**Theorem 10.4.** *If  $\gcd(a, b) = 1$  and  $a|z$  and  $b|z$  then  $ab|z$ .*

**Proof.** From the assumptions:  $z = ua = vb$ , therefore,  $a|vb$ . Since  $a, b$  are relative primes,  $a|v$ . (Here is a formal justification for such an observation: by 10.2,  $ax + by = 1$  for some  $x, y$ , hence  $vax + vby = v$ . Notice, that  $a$  divides both  $vax$  and  $vby$ , therefore,  $a|v$ .) Furthermore  $v = wa, z = vb = wab$  and  $ab|z$ .

The following generalization of 10.4 naturally holds: if  $m_1, m_2, \dots, m_n$  be pairwise relatively prime and  $m_i|z$  for all  $i = 1, 2, \dots, n$ , then  $m_1 \cdot m_2 \cdot \dots \cdot m_n|z$ .

Systems of linear congruences (consider a special case only): find  $x$  such that

$$x \equiv 2 \pmod{3}, \quad x \equiv 3 \pmod{5}, \quad x \equiv 1 \pmod{7}$$

**Theorem 10.5.** (The Chinese Remainder Theorem)

Let  $m_1, m_2, \dots, m_n$  be pairwise relatively prime. Then the system

$$x \equiv a_1 \pmod{m_1}$$

$$x \equiv a_2 \pmod{m_2}$$

.....

$$x \equiv a_n \pmod{m_n}$$

has a unique solution modulo  $m = m_1 \cdot m_2 \cdot \dots \cdot m_n$ .

**Proof.** For each  $k = 1, 2, \dots, n$  consider  $M_k = m/m_k = m_1 \cdot \dots \cdot m_{k-1} \cdot m_{k+1} \cdot \dots \cdot m_n$ . Note that  $\gcd(M_k, m_k) = 1$ , o.w. some  $d > 1$  divides both  $m_k$  and  $M_k$ , therefore  $d$  divides one of  $m_i$  for  $i \neq k$ , and  $m_i, m_k$  are not relatively prime. By 10.3,  $\exists y_k M_k y_k \equiv 1 \pmod{m_k}$ . Then  $x := a_1 M_1 y_1 + \dots + a_n M_n y_n$  is a desired solution. Indeed,  $m_i | M_j$  for all  $i \neq j$ , therefore  $x \equiv a_i M_i y_i \equiv a_i \cdot 1 \equiv a_i \pmod{m_i}$  for all  $i = 1, 2, \dots, n$ . Let us show the uniqueness. Suppose there is another nonnegative  $y < m$  such that  $y \equiv a_i \pmod{m_i}$ ,  $i = 1, 2, \dots, n$ . Without loss of generality assume that  $x \geq y$  and take the difference  $z = x - y$ . From the assumptions it follows that  $0 \leq z < m$  and  $z \equiv 0 \pmod{m_i}$ ,  $i = 1, 2, \dots, n$ . Therefore,  $m_i | z$  for all  $i = 1, 2, \dots, n$ . By 10.4 (the general form),  $m = m_1 \cdot m_2 \cdot \dots \cdot m_n | z$ , therefore,  $z = 0$ , i.e.  $x = y$ .

**Example 10.6.** To solve the system of congruences preceding 10.5, apply the general method from the proof of 10.5:  $m = 3 \cdot 5 \cdot 7 = 105$ ,  $M_1 = 5 \cdot 7 = 35$ ,  $M_2 = 3 \cdot 7 = 21$ ,  $M_3 = 3 \cdot 5 = 15$ ,  $35 \cdot 2 \equiv 1 \pmod{3}$ ,  $21 \cdot 1 \equiv 1 \pmod{5}$ ,  $15 \cdot 1 \equiv 1 \pmod{7}$ ,  $x = 2 \cdot 35 \cdot 2 + 3 \cdot 21 \cdot 1 + 2 \cdot 15 \cdot 1 = 233 \equiv 23 \pmod{105}$ .

**Example 10.7.** Handling large numbers by their remainders with respect to several smaller relative primes.  $m_1 = 99$ ,  $m_2 = 98$ ,  $m_3 = 97$ ,  $m_4 = 95$ ,  $m = m_1 \cdot m_2 \cdot m_3 \cdot m_4 = 89403930$ . Every  $k < m$  can be uniquely represented by a 4-tuple of numbers  $< 100$  that are the remainders of  $k$  with respect to  $m_1, m_2, m_3, m_4$ .  $123684 = (33, 8, 9, 89)$ ,  $413456 = (32, 92, 42, 16)$ . Therefore,  $123684 + 413456 = (65, 2, 51, 10)$ . To convert this 4-tuple back to the integer, one has to solve the system of congruences:  $x \equiv 65 \pmod{99}$ ,  $x \equiv 2 \pmod{98}$ ,  $x \equiv 51 \pmod{97}$ , and  $x \equiv 10 \pmod{99}$ .

**Theorem 10.8.** (Fermat's Little Theorem)

If  $p$  is a prime which does not divide  $a$  then  $a^{p-1} \equiv 1 \pmod{p}$ . Furthermore,  $a^p \equiv a \pmod{p}$ .

**Proof.** Consider  $\mathbf{Z}_p^+ = \{1, 2, 3, \dots, p-1\}$  the set of all positive remainders modulo  $p$ , and let  $a\mathbf{Z}_p^+ = \{a \cdot 1, a \cdot 2, a \cdot 3, \dots, a \cdot (p-1)\}$ . All elements in the latter set are distinct  $\pmod{p}$ . Indeed, let  $ax \equiv ay \pmod{p}$  and  $x \geq y$ , thus  $0 \leq (x - y) < p$ . Then  $a(x - y) \equiv 0 \pmod{p} \Rightarrow p | a(x - y) \Rightarrow p | a$  or  $p | (x - y)$ . The former is impossible by the assumptions of the theorem. Therefore,  $p | (x - y)$  and thus  $x - y = 0$ , i.e.  $x = y$ . We have established, that  $\mathbf{Z}_p^+$  and  $a\mathbf{Z}_p^+$  is the same set modulo  $p$ , therefore, the products of their elements coincide mod  $p$ :

$$1 \cdot 2 \cdot \dots \cdot (p-1) \equiv (a \cdot 1) \cdot (a \cdot 2) \cdot \dots \cdot (a \cdot (p-1)) \pmod{p}$$

$$(p-1)! \equiv a^{p-1} (p-1)! \pmod{p}, \quad (p-1)! (a^{p-1} - 1) \equiv 0 \pmod{p}, \text{ thus } p | (p-1)! \text{ or } p | (a^{p-1} - 1).$$

The former is impossible since a prime  $p$  cannot divide any positive number  $< (p-1)$ . Therefore  $p | (a^{p-1} - 1)$  and  $a^{p-1} \equiv 1 \pmod{p}$ .

**Example 10.9.** Evaluate  $2^{340} \pmod{11}$ . By Fermat's Little Theorem,  $2^{10} \equiv 1 \pmod{11}$ . Therefore  $2^{340} = (2^{10})^{34} \equiv 1^{34} \equiv 1 \pmod{11}$ .

**RSA encryption.** See the slides and/or the textbook.

**Homework assignments.** (due Friday 02/16).

10A:Rosen2.5-2f; 10B:Rosen2.5-24ab; 10C:Rosen2.5-26di