### **Priority Queues and Heaps**

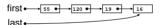
# The Bag Interface

```
interface Bag<E> {
    void put(E obj);
    E get(); //extract some element
    boolean isEmpty();
}
```

Examples: Stack, Queue

### Stacks and Queues as Lists

- Stack (LIFO) implemented as list
   put(), get() from front of list
- Queue (FIFO) implemented as list
- put() on back of list, get() from front of list
- All Bag operations are O(1)



### **Priority Queue**

- A Bag in which data items are Comparable
- lesser elements (as determined by compareTo()) have higher priority
- get() returns the element with the highest priority = least in the compareTo() ordering
- break ties arbitrarily

# **Examples**

- Scheduling jobs to run on a computer
  - default priority = arrival time
  - priority can be changed by operator
- Scheduling events to be processed by an event handler
  - priority = time of occurrence
- · Airline check-in
  - first class, business class, coach
  - FIFO within each class

### **Priority Queues**

```
interface Bag<E> {
   void put(E obj);
   E get(); //extract some element
   boolean isEmpty();
}
interface PriorityQueue<E extends Comparable>
   extends Bag<E> {}
```

### **Priority Queues as Lists**

- · Maintain as unordered list
  - put ( ) puts new element at front O(1)
  - get() must search the list O(n)
- · Maintain as ordered list
  - put() must search the list O(n)
  - get() gets element at front O(1)
- In either case, O(n2) to process n elements

Can we do better?

### Important Special Case

- Fixed number of priority levels 0,...,p − 1
- FIFO within each level
- Example: airline check-in
- put ( ) insert in appropriate queue O(1)
- •get() must find a nonempty queue O(p)

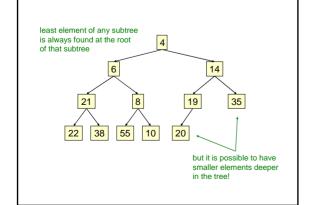
### Heaps

- A *heap* is a concrete data structure that can be used to implement priority queues
- Gives better complexity than either ordered or unordered list implementation:
  - put(), get() O(log n)
     isEmpty() O(1)
- O(n log n) to process n elements
- Do not confuse with heap memory, where the Java virtual machine allocates space for objects - different usage!

### Heaps

- Binary tree with data at each node
- Satisfies the Heap Order Invariant.

The least (highest priority) element of any subtree is found at the root of that subtree



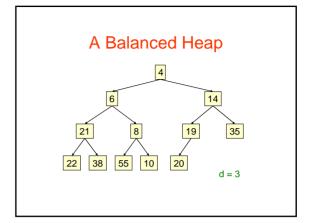
# **Examples of Heaps**

- Ages of people in family tree
  - parent is always older than children, but you can have an uncle who is younger than you
- Salaries of employees of a company
  - bosses generally make more than subordinates, but a VP in one subdivision may make less than a Project Supervisor in a different subdivision

### **Balanced Heaps**

### Two restrictions:

- 1. Any node of depth < d 1 has exactly 2 children, where d is the height of the tree
  - implies that any two maximal paths (path from a root to a leaf) are of length d or d – 1, and the tree has at least 2<sup>d</sup> nodes
- 2. All maximal paths of length d are to the left of those of length d 1



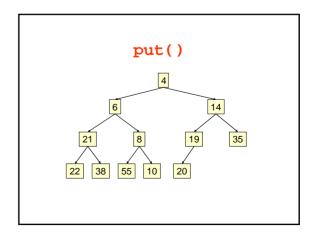
## Store in an Array or Vector

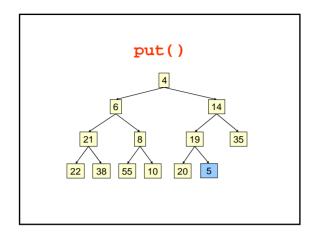
- Elements of the heap are stored in the array in order, going across each level from left to right, top to bottom
- The children of the node at array index n are found at 2n + 1 and 2n + 2
- The parent of node n is found at (n-1)/2

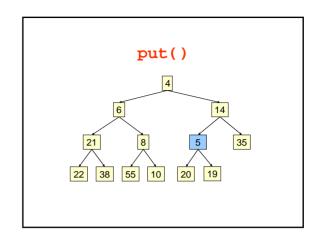
# Store in an Array or Vector 40 114 21 38 55 10 10 20 11 children of node n are found at 2n + 1 and 2n + 2

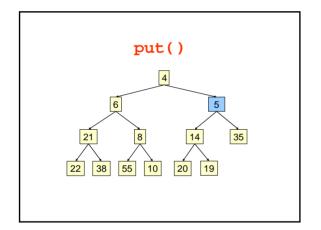
### put()

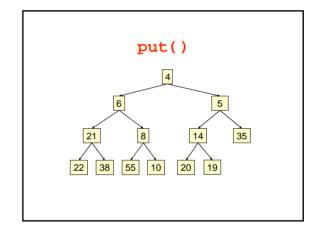
- Put the new element at the end of the array
- If this violates heap order because it is smaller than its parent, swap it with its parent
- Continue swapping it up until it finds its rightful place
- The heap invariant is maintained!

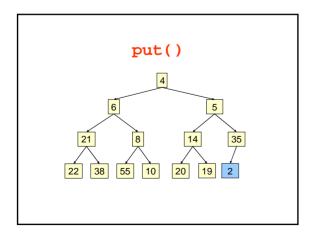


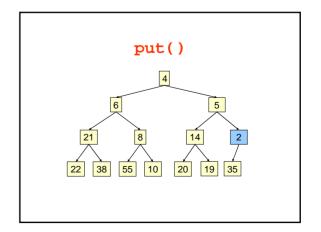


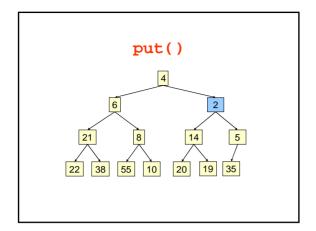


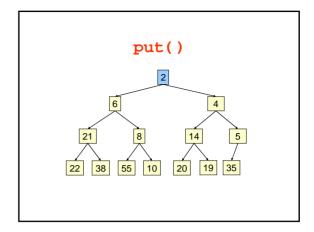


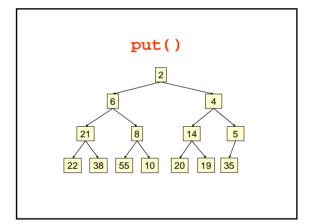












# put()

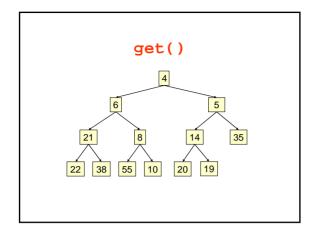
Time is O(log n), since the tree is balanced

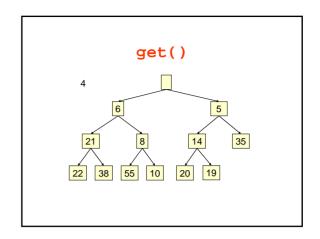
 size of tree is exponential as a function of depth
 depth of tree is logarithmic as a function of size

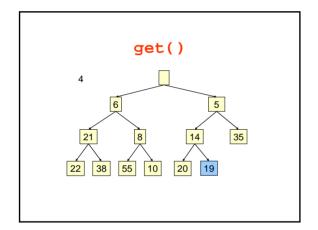
```
class PQ<E extends Comparable> extends java.util.Vector<E>
implements PriorityQueue<E> {
  public void put(E obj) {
    add(obj); //add new element to end of array
    rotateUp(size() - 1);
  }
  private void rotateUp(int index) {
    if (index == 0) return;
    int parent = (index - 1)/2;
    if (elementAt(parent).compareTo(elementAt(index)) <= 0)
        return;
    swap(index, parent);
    rotateUp(parent);
  }
}</pre>
```

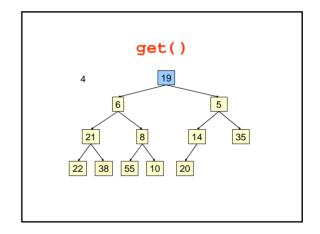
### get()

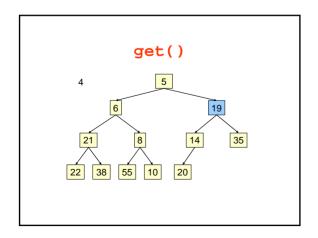
- Remove the least element it is at the root
- This leaves a hole at the root fill it in with the last element of the array
- If this violates heap order because the root element is too big, swap it down with the smaller of its children
- Continue swapping it down until it finds its rightful place
- The heap invariant is maintained!

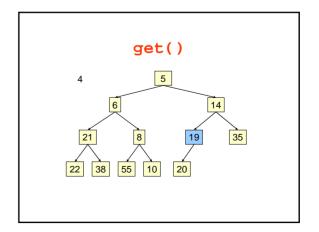


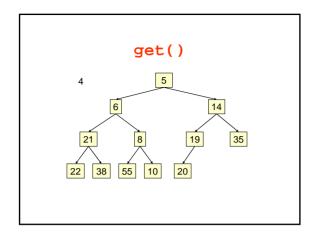


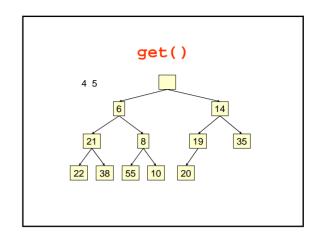


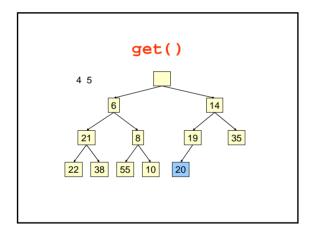


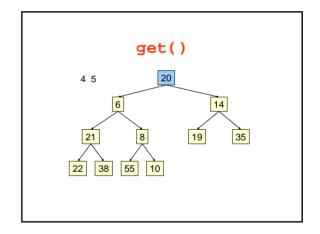


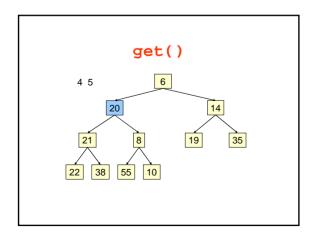


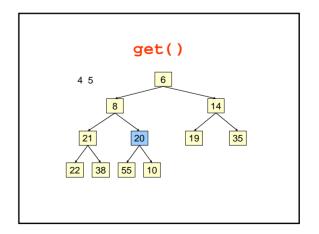


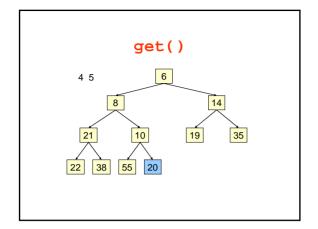


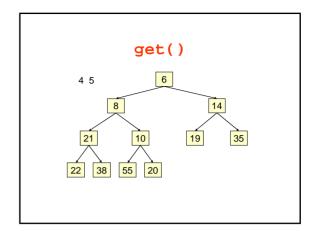












### get()

• Time is O(log n), since the tree is balanced

```
public E get() {
   if (isEmpty()) throw new NoSuchElementException();
   E temp = elementAt(0);
   setElementAt(elementAt(size() - 1), 0);
   setSize(size() - 1);
   rotateDown(0);
   return temp;
}
private void rotateDown(int index) {
   int child = 2*(index + 1); //right child
   if (child >= size()
   || elementAt(child - 1).compareTo(elementAt(child)) < 0)
        child -= 1;
   if (child >= size()) return;
   if (elementAt(index).compareTo(elementAt(child)) <= 0)
        return;
   swap(index, child);
   rotateDown(child);
}</pre>
```

### **HeapSort**

Given a Comparable[] array of length n,

- 1. Put all n elements into a heap O(n log n)
- 2. Repeatedly get the min O(n log n)

```
public static void heapSort(Comparable[] a) {
   PriorityQueue<Comparable> pq = new PQ<Comparable>();
   for (Comparable x : a) { pq.put(x); }
   for (int i = 0; i < a.length; i++) { a[i] = pq.get(); }
}</pre>
```