Previous Lecture:

- Vectorized operations
- Introduction to 2-d array—matrix

Today, Lecture 14:

- Examples on computing with matrices
- Pre-reading: 7.1 (transition probability matrices)
- Post-reading: contour plots, PDEs (7.2, 7.3 in *Insight*)

Announcements:

- Prelim 1 tonight at 6:30pm
 - Barton Hall, west entrance; see CMS for seat assignment
 - Online: see Canvas for Zoom, Gradescope links
- Project 4 posted by Thurs, due April 8

Pattern for traversing a matrix M

```
[nr, nc] = size(M)
for r= I:nr
   % At row r
    for c= 1:nc
         % At column c (in row r)
         % Do something with M(r,c) ...
    end
end
```

Exercise: what's different about this version?

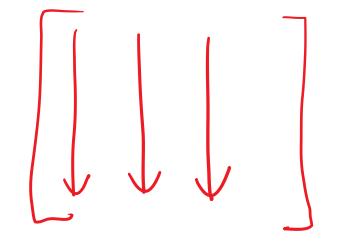
```
function val = minInMatrix(M)
val = M(1,1);
[nr, nc] = size(M);
for c = 1:nc
    for r = 1:nr
        if M(r,c) < val
            val= M(r,c);
        end
    end
end
```

A: Nothing (identical to previous version)

B: Searches elements in a different order

C: Index-out-of-bounds if M is not square

D: Doesn't correctly return minimum





A Cost/Inventory Problem

- A merchant has 3 supplier warehouses that stock 5 different products
- The cost of a product varies from warehouse to warehouse
- The inventory varies from warehouse to warehouse

Problems

A customer submits a purchase order that is to be shipped from a single warehouse.

- I. How much would it cost a warehouse to fill the order?
- 2. Does a warehouse have enough inventory to fill the order?
- 3. Among the warehouses that can fill the order, who can do it most cheaply?

Available data

C(i,j) is what it costs warehouse i to supply product j

Inv(i,j) is the
 inventory in
 warehouse i of
 product j

Inv

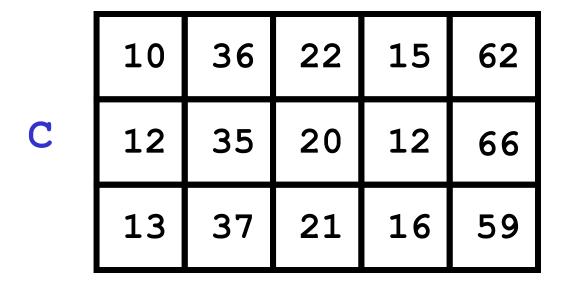
PO(j) is the number of product j's that the client wants

10	36	22	15	62
12	35	20	12	66
13	37	21	16	59

 38
 5
 99
 34
 42

 82
 19
 83
 12
 42

 51
 29
 21
 56
 87

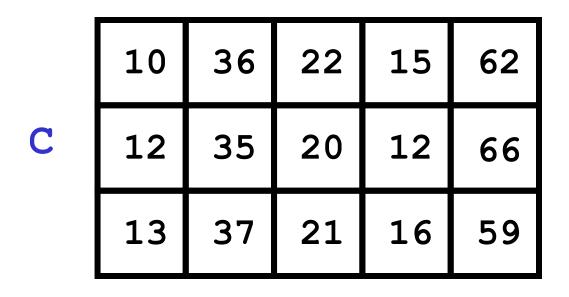


PO

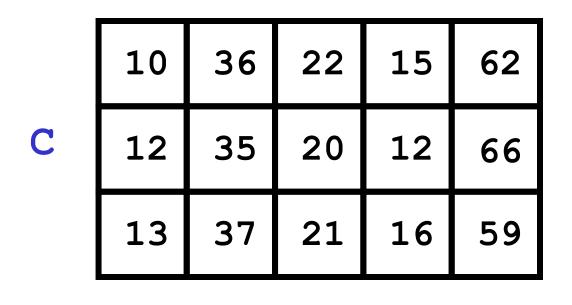
1	0	12	29	5

example out specific generalize

Cost for warehouse 1:



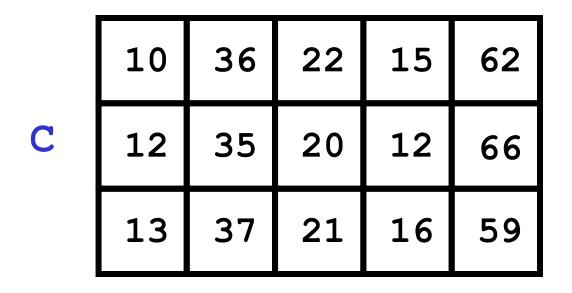
```
Cost for s = 0; %Sum of cost warehouse 1: for j=1:5 % n cols s = s + C(1,j) *PO(j) end
```



```
Cost for s = 0; %Sum of cost warehouse 2: for j=1:5
```

$$s = s + C(2,j)*PO(j)$$

end



```
Cost for s = 0; %Sum of cost warehouse i: for j=1:5
s = s + C(i,j)*PO(j) end
```

```
function theBill = iCost(i,C,PO)
% The cost when warehouse i fills the
% purchase order
nProd= length(PO); or Size(c,2)
theBill= 0;
for j = 1:nProd
   theBill + C(i,j)*PO(j);
end
```

Finding the Cheapest

Both which warehouse and how cheap

```
iBest= 0; minBill=(inf)
for i = 1:nSuppliers
   iBill= iCost(i,C,PO);
   if iBill < minBill
      % Found an Improvement
      iBest= i; minBill= iBill;
   end
end
```

Aside: floating-point "bonus numbers"

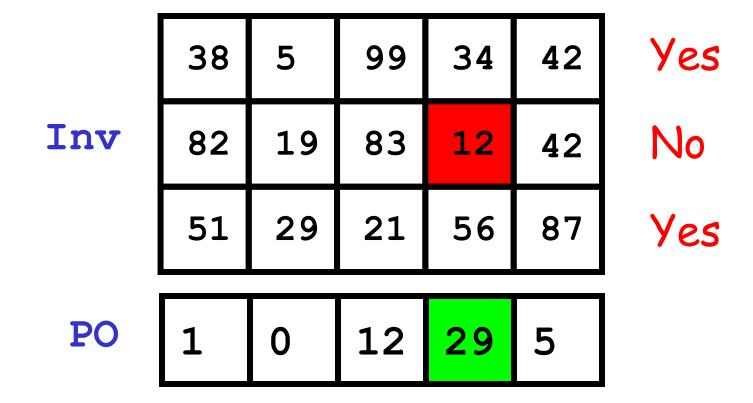
- inf: Represents "infinity"
 - Both positive and negative versions
 - Larger (or smaller) than any other number
 - Generated on overflow or when dividing by zero
- nan: Not-a-number
 - Not equal to anything (even itself)
 - Not greater-than or less-than anything (even inf)
 - Generated from 0/0, inf*0, ...
 - If involved in mathematical operation, result will be nan

Inventory/Capacity Considerations

What if a warehouse lacks the inventory to fill the purchase order?

Such a warehouse should be excluded from the find-the-cheapest computation.

Who Can Fill the Order?



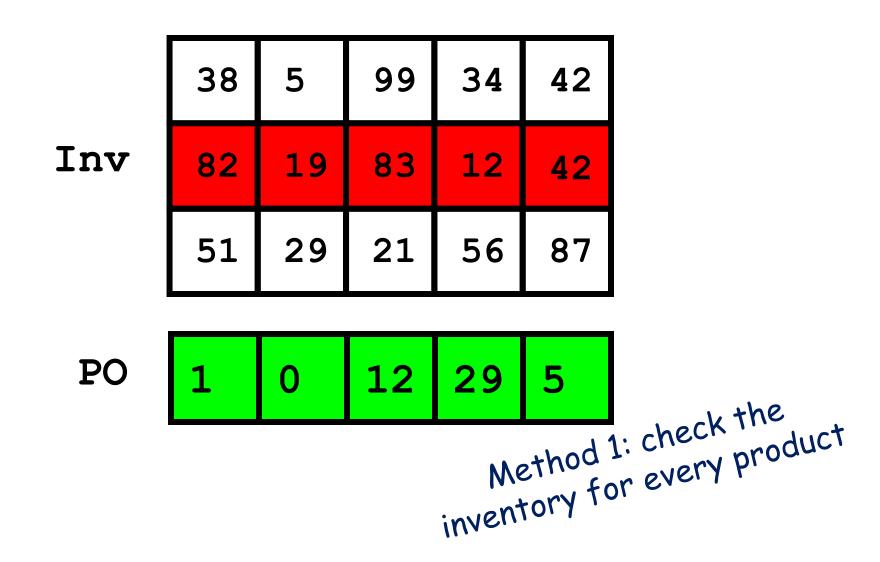
Wanted: A True/False Function



DO is "true" if warehouse i can fill the order.

DO is "false" if warehouse i cannot fill the order.

Example: Check inventory of warehouse 2



Initialization

	38	5	99	34	42		
Inv	82	19	83	12	42	DO	1
	51	29	21	56	87		
PO	1	0	12	29	5		

Still True...

	38	5	99	34	42		
Inv	82	19	83	12	42	DO	1
	51	29	21	56	87		
PO	1	0	12	29	5		

$$DO = (Inv(2,1) >= PO(1))$$

Still True...

	38	5	99	34	42		
Inv	82	19	83	12	42	DO	1
	51	29	21	56	87		
PO	1	0	12	29	5		

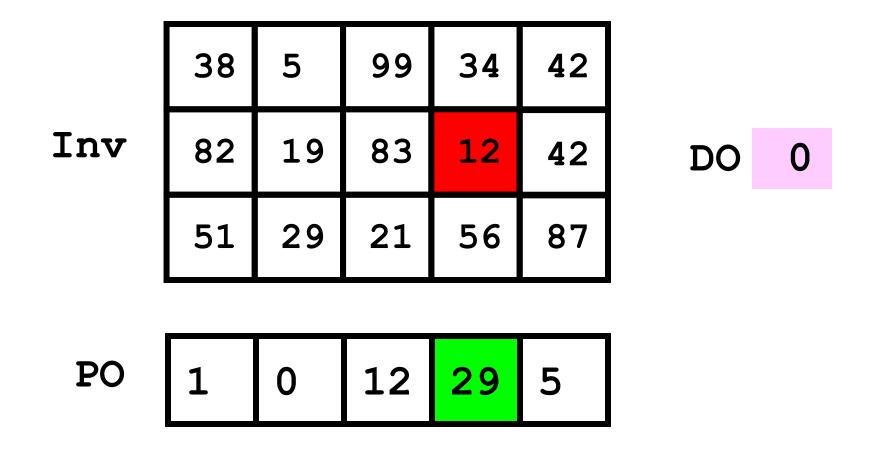
$$DO = DO \&\& (Inv(2,2) >= PO(2))$$

Still True...

	38	5	99	34	42		
Inv	82	19	83	12	42	DO	1
	51	29	21	56	87		
PO	1	0	12	29	5		

$$DO = DO && (Inv(2,3) >= PO(3))$$

No Longer True...



$$DO = DO \&\& (Inv(2,4) >= PO(4))$$

Stay False...

	38	5	99	34	42		
Inv	82	19	83	12	42	DO	0
	51	29	21	56	87		
PO	1	0	12	29	5		

$$DO = DO && (Inv(2,5) >= PO(5))$$

```
function DO = iCanDo(i,Inv,PO)
% DO is true if warehouse i can fill
% the purchase order. Otherwise, false
nProd= length(PO);
DO=1;
for j = 1:nProd
     DO= DO && ( Inv(i,j) >= PO(j) );
end
```

```
function DO = iCanDo(i,Inv,PO)
% DO is true if warehouse i can fill
% the purchase order. Otherwise, false
nProd= length(PO);
j=1;
while j \le nProd && Inv(i,j) >= PO(j)
   j = j + 1;
                    Method 2: stop as soon as you find one product for which there isn't enough inventory
end
DO=
```

```
iClicker+

FOWER

B

COMBANIES

A

B

C

D

B

E
```

```
function DO = iCanDo(i,Inv,PO)
% DO is true if warehouse i can fill
% the purchase order. Otherwise, false
nProd= length(PO);
j=1;
while j \le nProd && Inv(i,j) >= PO(j)
   j = j + 1;
                  DO should be true when...
end
                        < nProd
                       == nProd
DO=
                      j > nProd
```

```
function DO = iCanDo(i,Inv,PO)
% DO is true if warehouse i can fill
% the purchase order. Otherwise, false
nProd= length(PO);
j=1;
while j \le nProd && Inv(i,j) >= PO(j)
   j = j + 1;
                          inv
end
DO= (j>nProd);
                           PO
                                 6
```

Back To Finding the Cheapest

```
iBest= 0; minBill= inf;
for i = 1:nSuppliers
   iBill= iCost(i,C,PO);
   if iBill < minBill</pre>
       % Found an Improvement
       iBest= i; minBill= iBill;
   end
               Don't bother with this unless there is sufficient inventory.
end
```

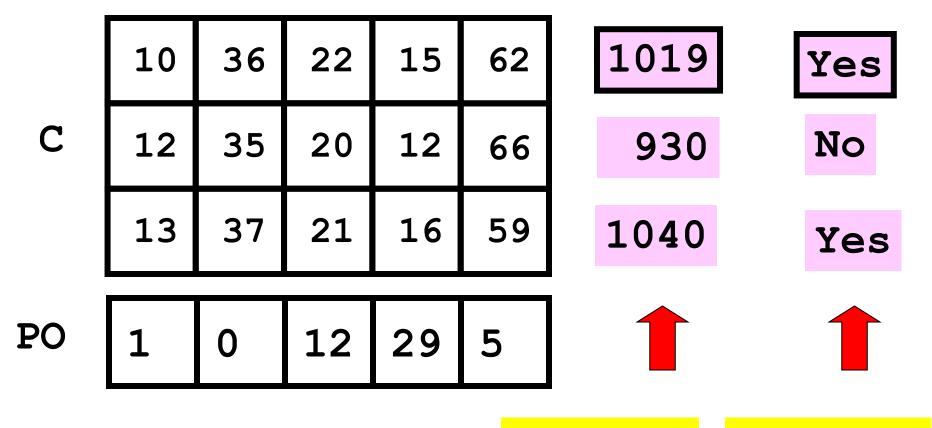
Back To Finding the Cheapest

```
iBest= 0; minBill= inf;
for i = 1:nSuppliers
   if iCanDo(i,Inv,PO)
      iBill= iCost(i,C,PO);
      if iBill < minBill
      % Found an Improvement
         iBest= i; minBill= iBill;
      end
   end
```

end

See Cheapest.m for alternative implementation

Finding the Cheapest



As computed by iCost

As computed by iCanDo

Matrix example: Random Web

N web pages can be represented by an N-by-N Link Array A.



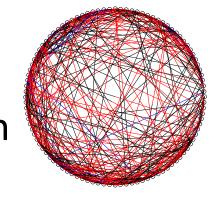
webpage i

Page 4 links to

0 0 1 0/1 0 0 Page 2 But Page 2 does 1011010 not link buck 0011011011 to Page 4 0110101

Matrix example: Random Web

 N web pages can be represented by an N-by-N Link Array A.



A(i,j) is 1 if there is a link on webpage j to webpage i

Generate a random link array and display the connectivity:

- There is no link from a page to itself
- If $i\neq j$ then A(i,j) = 1 with probability $\frac{1}{1+|i-j|}$ There is more likely to be a link if i is close to j

```
function A = RandomLinks(n)
% A is n-by-n matrix of 1s and 0s
% representing n webpages
A= zeros(n,n); % initialize to 0s
for i = 1:n
  for j = 1:n
    % if A(i,j) not on diagonal,
    % assign 1 with some probability
  end
end
```

An event happens with probability p, 0<p<1

```
% Flip a fair coin
r= rand();
if r < .5
    disp('heads')
else
    disp('tails')
end</pre>
```

```
% Unfair coin: shows heads
% twice as often as tails
r= rand();
if r < 2/3
    disp('heads')
else
    disp('tails')
end</pre>
```

```
% Event X happens with probability p
r= rand();
if r
```

```
function A = RandomLinks(n)
% A is n-by-n matrix of 1s and 0s
% representing n webpages
A = zeros(n,n);
for i = 1:n
  for j = 1:n
   r= rand();
    if i~=j && r < 1/(1 + abs(i-j));
        A(i,j) = 1;
    end
  end
end
```