Lecture 26

Advanced Sorting

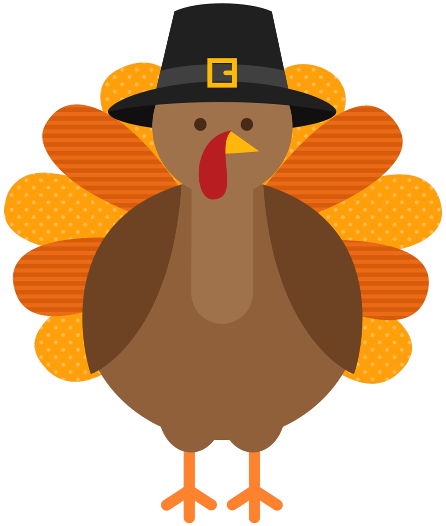
Announcements for This Lecture

Exam & Assignments

- **Prelim, Dec. 5 at 7:30**
  - See webpage for rooms
  - Last call for conflicts!
- Review **Dec. 4** at 5pm
- A6 is now graded
  - **Mean**: 92.5 **Median**: 96.5
  - **Time**: 12.3hrs **St Dev**: 6.4hr
- Keep working on A7
  - Labs today are open OH

Optional Videos

- **ALL** all are now posted
  - **Lesson 28** for today
  - **Lesson 29** is next week



Recall Our Problem

- Both insertion, selection sort are **nested loops**
  - **Outer loop** over each element to sort
  - **Inner loop** to put next element in place
  - Each loop is n steps. $n \times n = n^2$
- To do better we must *eliminate* a loop
  - But how do we do that?
  - What is like a loop? **Recursion!**
  - First need an *intermediate* algorithm

The Partition Algorithm

- Given a list segment b[h..k] with some value x in b[h]:

Start: b

h	k
x	?

- Swap elements of b[h..k] to get this answer

Goal: b

h	i	i+1 k
<= x	x	>= x

change: b

h								k
3	5	4	1	6	2	3	8	1

into b

h			i					k
1	2	1	**3**	5	4	6	3	8

or b

h				i				k
1	2	3	1	**3**	4	5	6	8

- x is called the pivot value
  - x is not a program variable
  - denotes value initially in b[h]

Designing the Partition Algorithm

- Given a list b[h..k] with some value x in b[h]:

Start: b

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Goal: b

h	i	i+1 k
<= **x**	**x**	>= **x**

In-Progress: b

h	i		j k
<= **x**	**x**	?	>= **x**

Indices b, h important!
Might partition only part

Implementating the Partition Algorithm

```
def partition(b, h, k):
    """Partition list b[h..k] around a pivot x = b[h]"""
    i = h; j = k+1; x = b[h]

    while i < j-1:
        if b[i+1] >= x:
            # Move to end of block.
            swap(b,i+1,j-1)
            j = j - 1
        else:   # b[i+1] < x
            swap(b,i,i+1)
            i = i + 1

    return i
```

partition(b,h,k), not partition(b[h:k+1])

Remember, slicing always copies the list!

We want to partition the **original** list

Partition Algorithm Implementation

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<= **x**	**x**	?	>= **x**
h	i	i+1	j k
1 2	**3**	1 5 0	6 3 8

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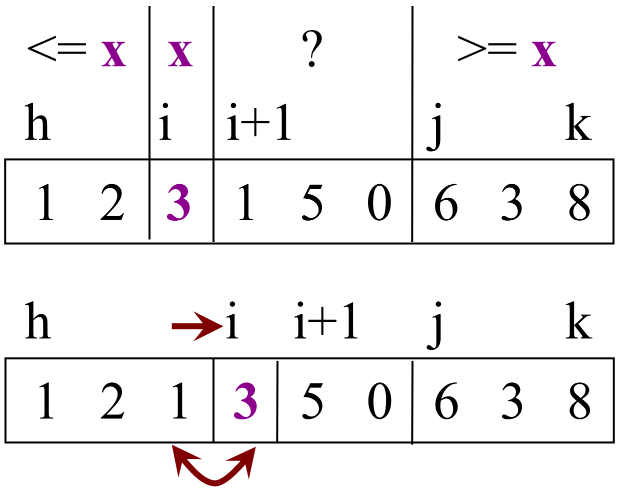

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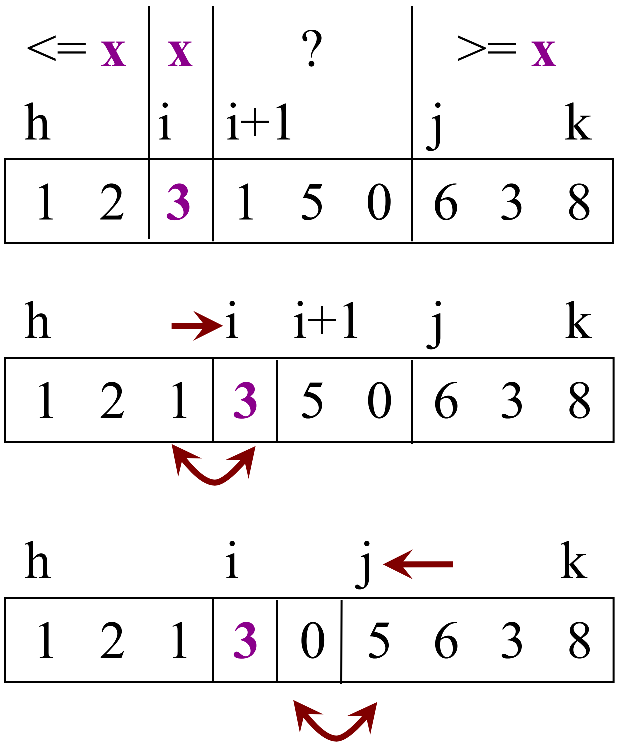

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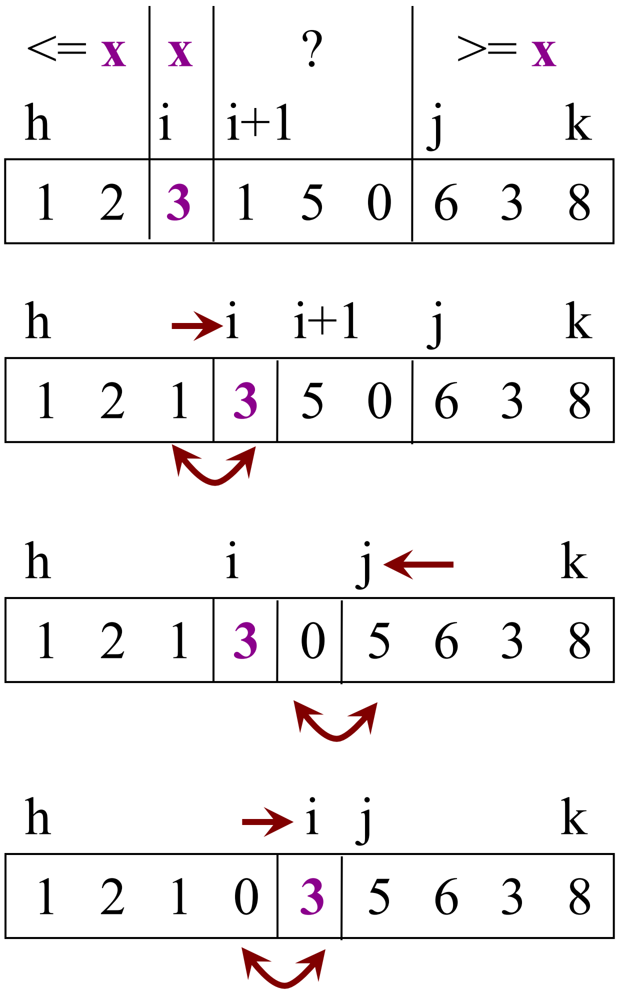
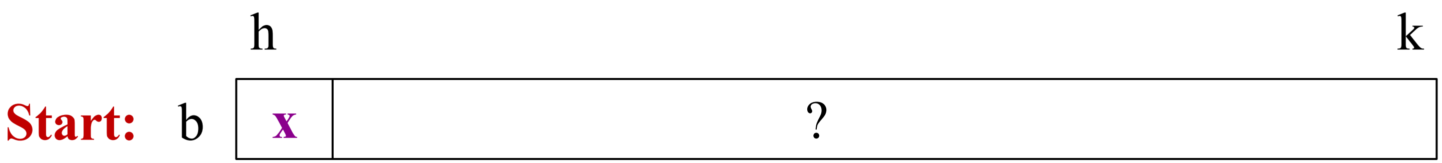


Why is this Useful?

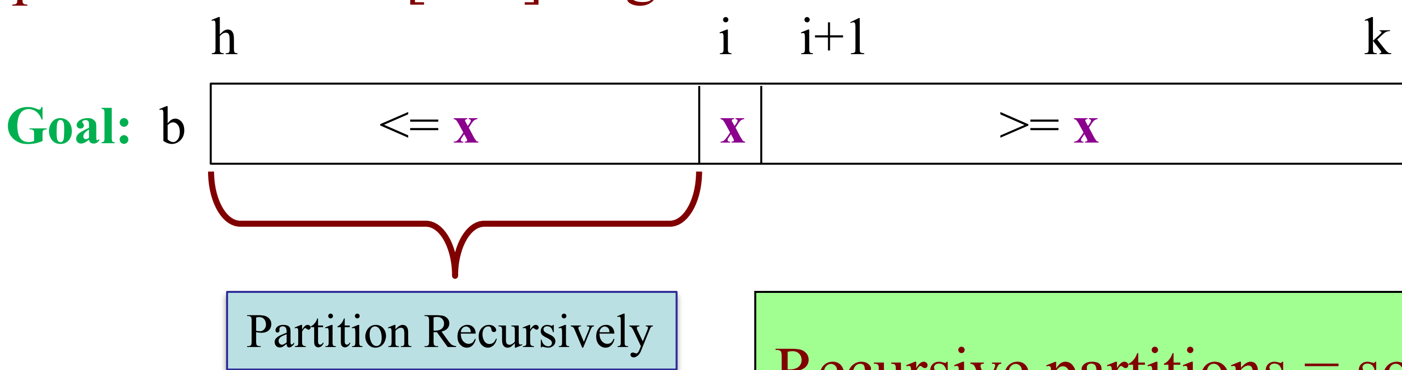
- Will use this algorithm to replace inner loop
  - The inner loop cost us n swaps every time
- Can this reduce the number of swaps?
  - Worst case is k-h swaps
  - This is n if partitioning the whole list
  - But less if only partitioning part
- **Idea:** Break up list and partition only part?
  - This is **Divide-and-Conquer!**

Sorting with Partitions

- Given a list segment b[h..k] with some value x in b[h]:



- Swap elements of b[h..k] to get this answer



Recursive partitions = sorting

- Called **QuickSort** (why???)
- Popular, fast sorting technique

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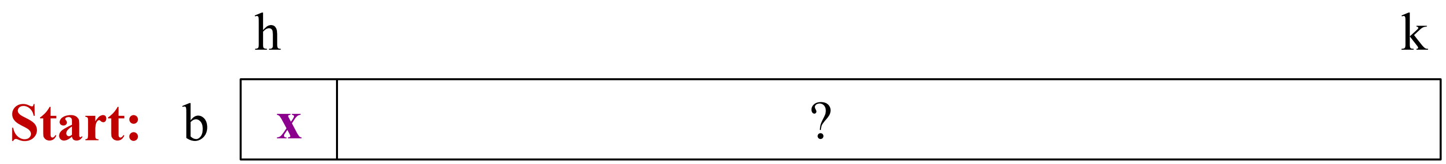

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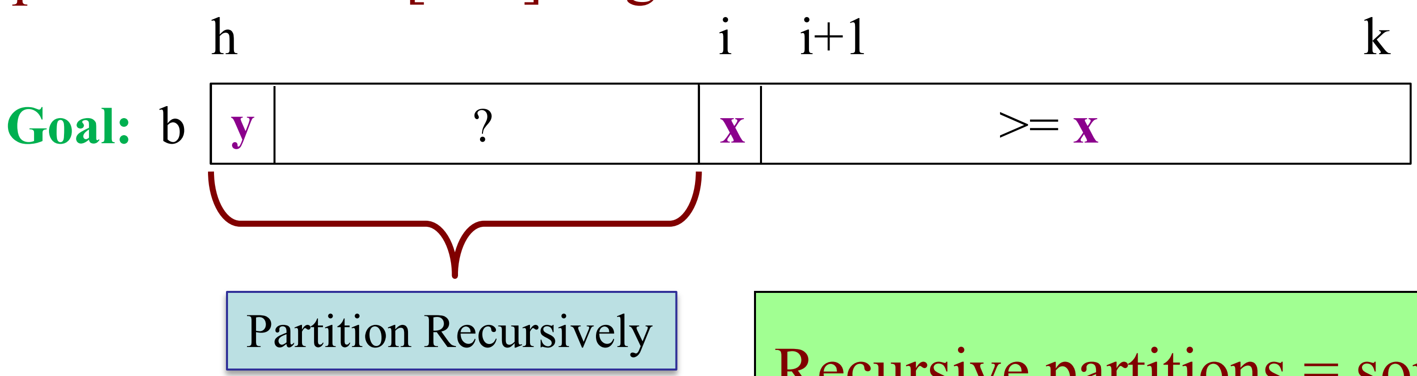

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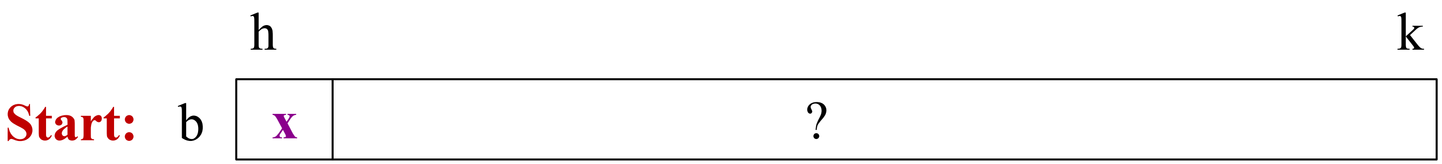

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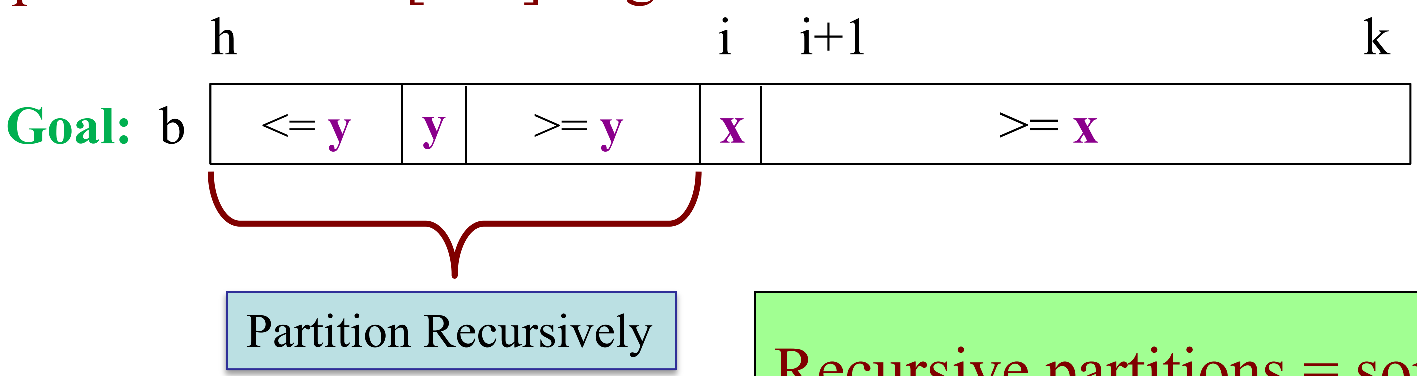

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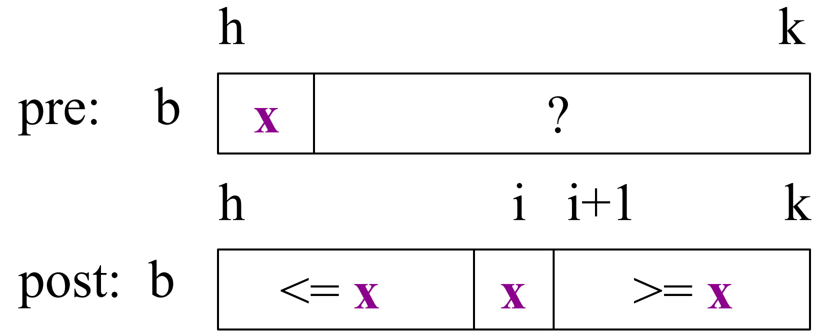
QuickSort

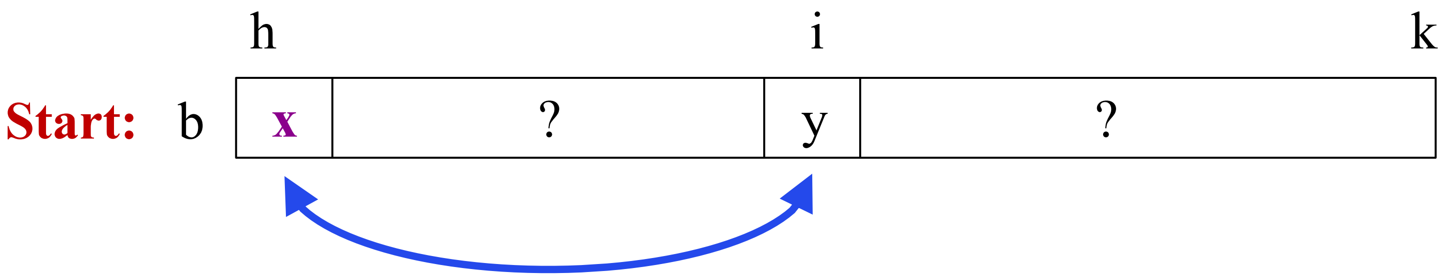
```
def quick_sort(b, h, k):
    """Sort the array fragment b[h..k]"""
    if  b[h..k] has fewer than 2 elements:
        return
    j = partition(b, h, k)
    # b[h..j–1] <= b[j] <= b[j+1..k]
    # Sort b[h..j–1]  and  b[j+1..k]
    quick_sort (b, h, j–1)
    quick_sort (b, j+1, k)
```

- **Worst Case:** array already sorted
  - Or almost sorted
  - n^2 in that case
- **Average Case:** array is scrambled
  - n log n in that case
  - Best sorting time!



So Does that Solve It?

- Worst case still seems bad! Still n^2
  - But only happens in small number of cases
  - Just happens that case is common (already sorted)
- Can greatly reduce issue with randomization
  - Swap start with random element in list
  - Now pivot is random and already sorted unlikely



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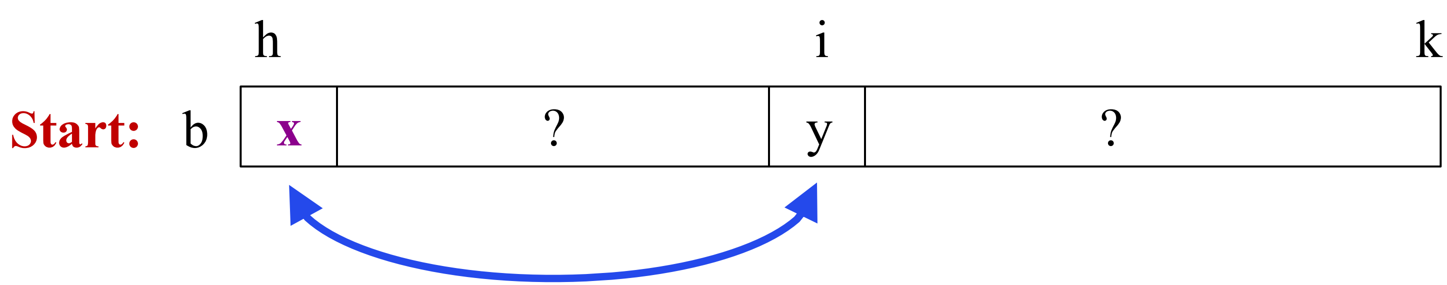
Makes it "good enough" for most applications



Can We Do Better?

- Recursion seems to be the solution
  - Partitioned the list into two halves
  - Recursively sorted each half
- How about a traditional **divide-and-conquer**?
  - **Divide** the list into two halves
  - **Recursively sort** the two halves
  - **Combine** the two sort halves
- How do we do the last step?

Combining Two Sorted Lists

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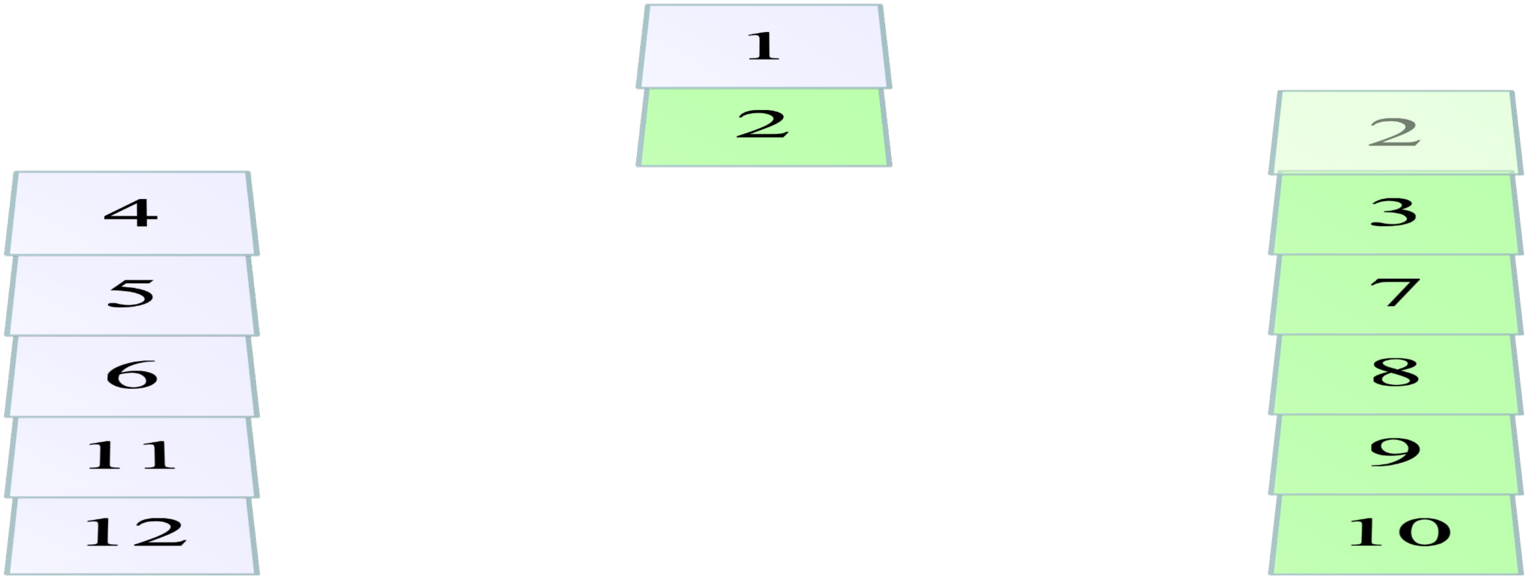
Pick from list with the least

Combining Two Sorted Lists



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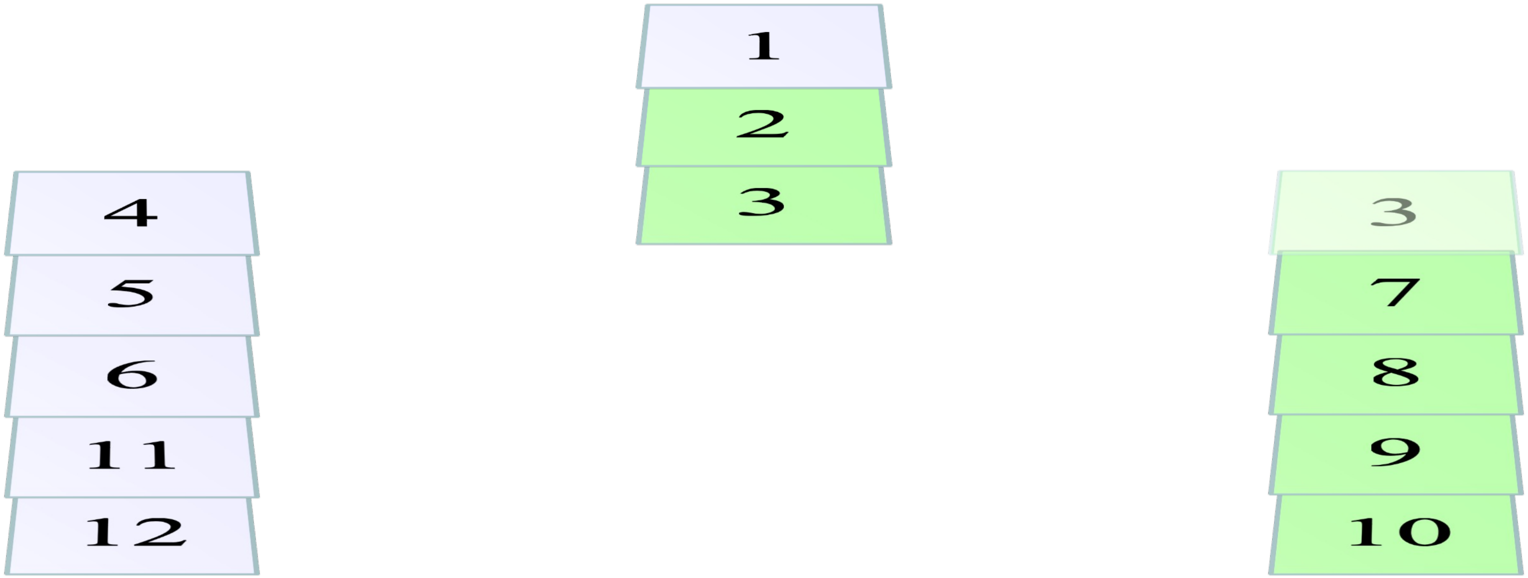

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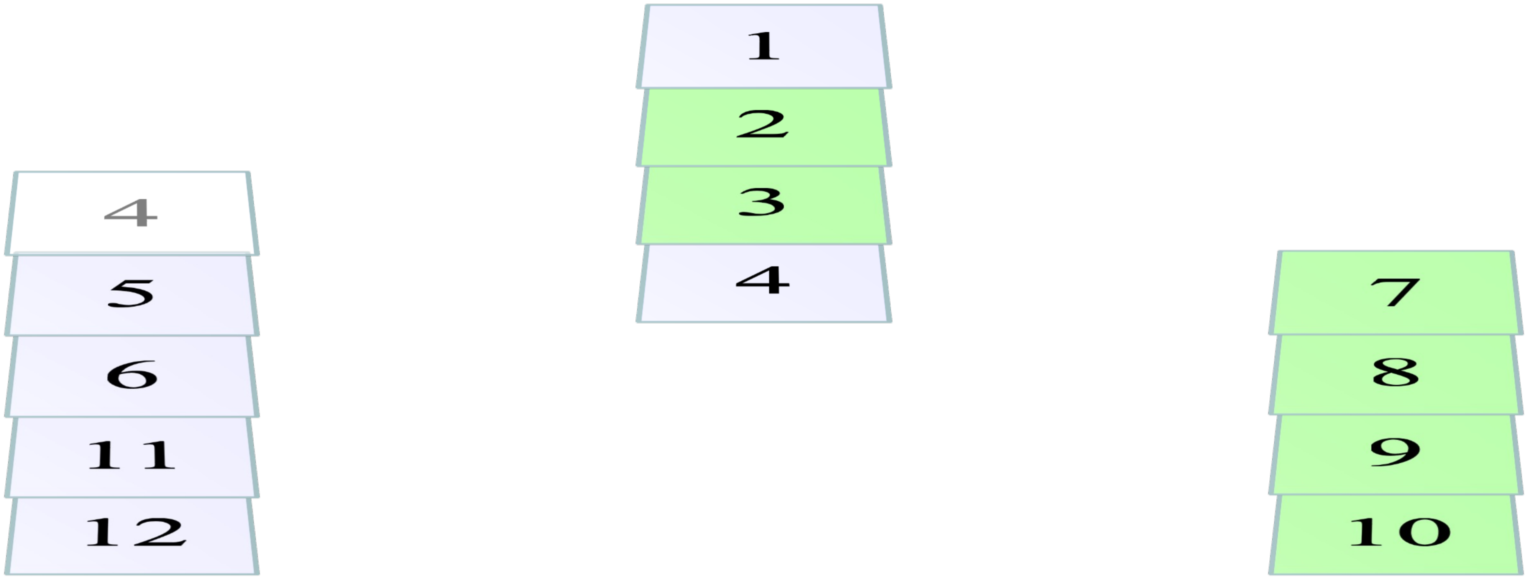

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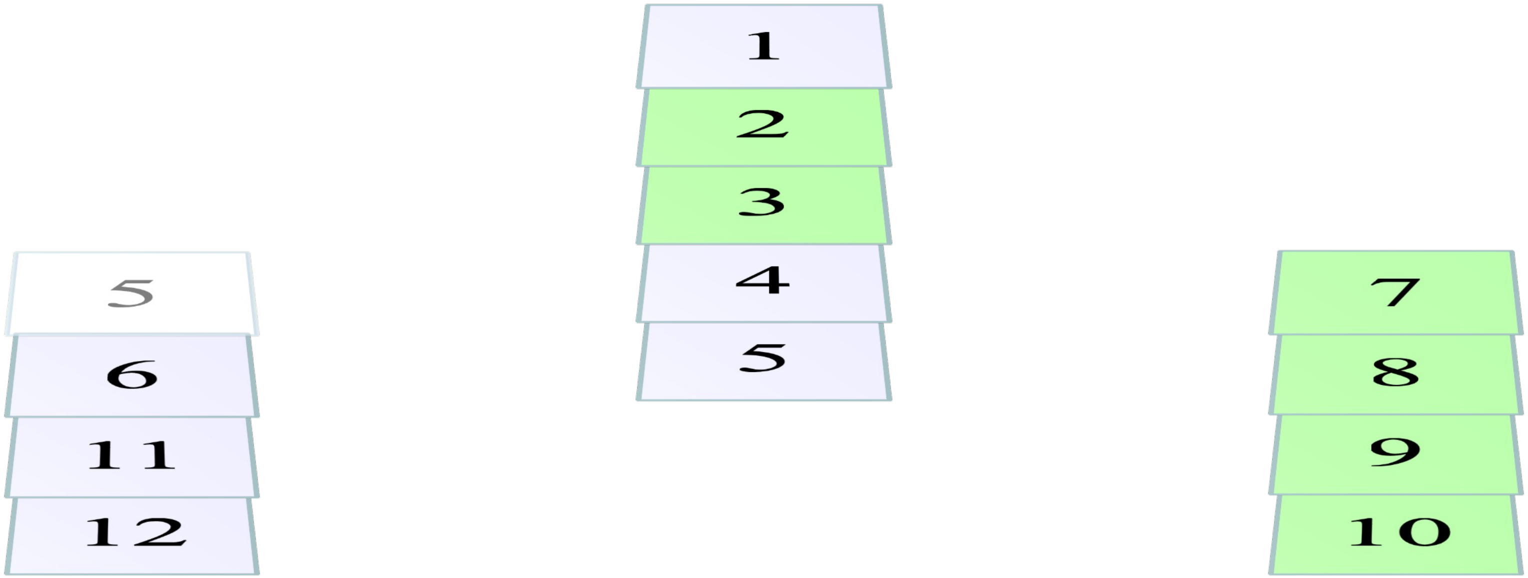

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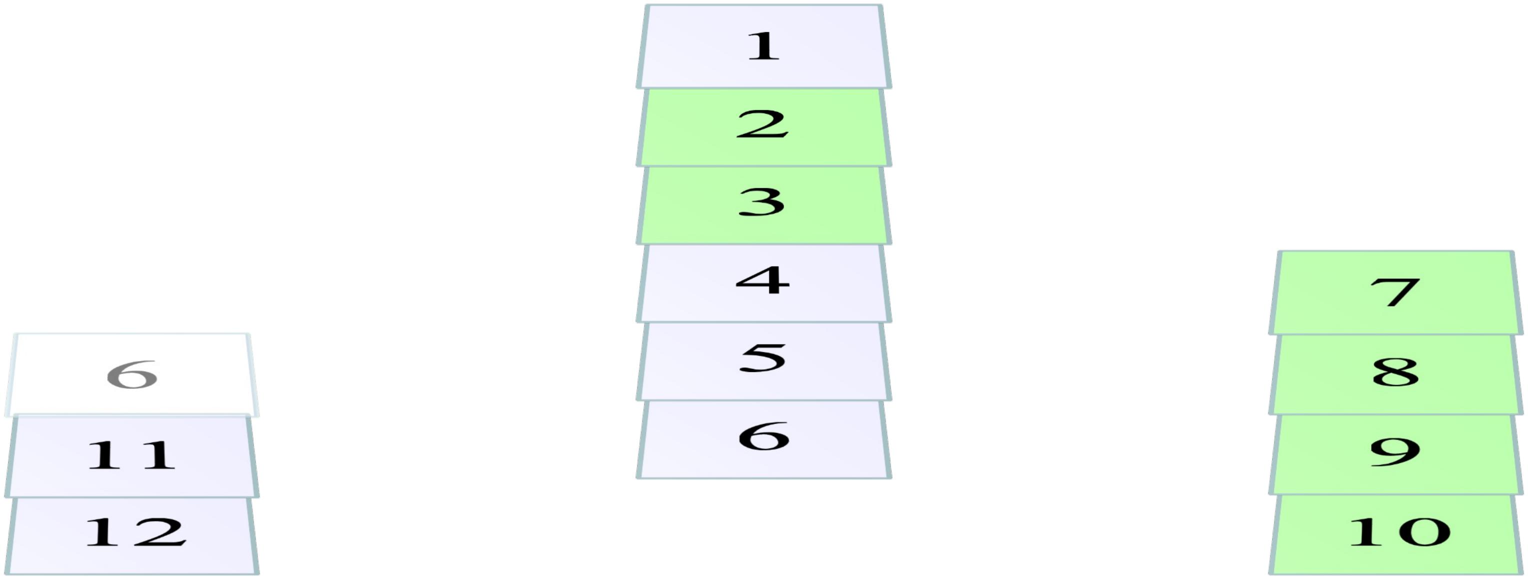

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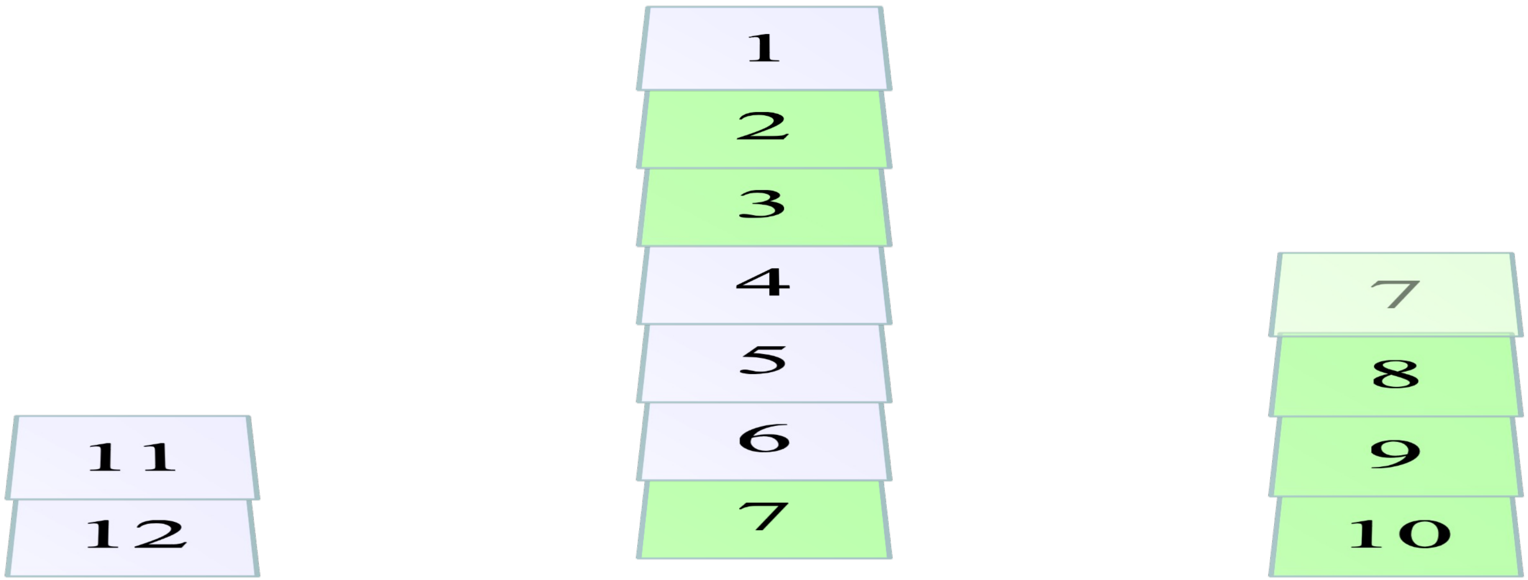

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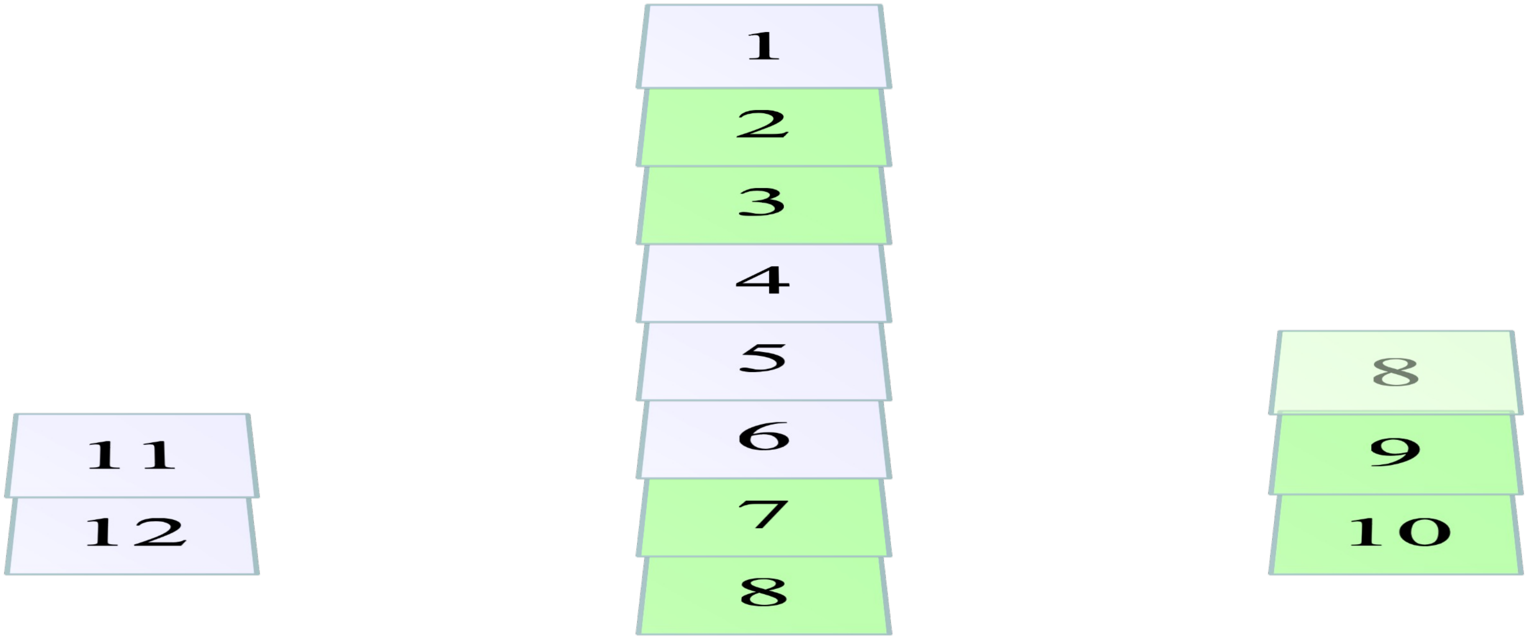

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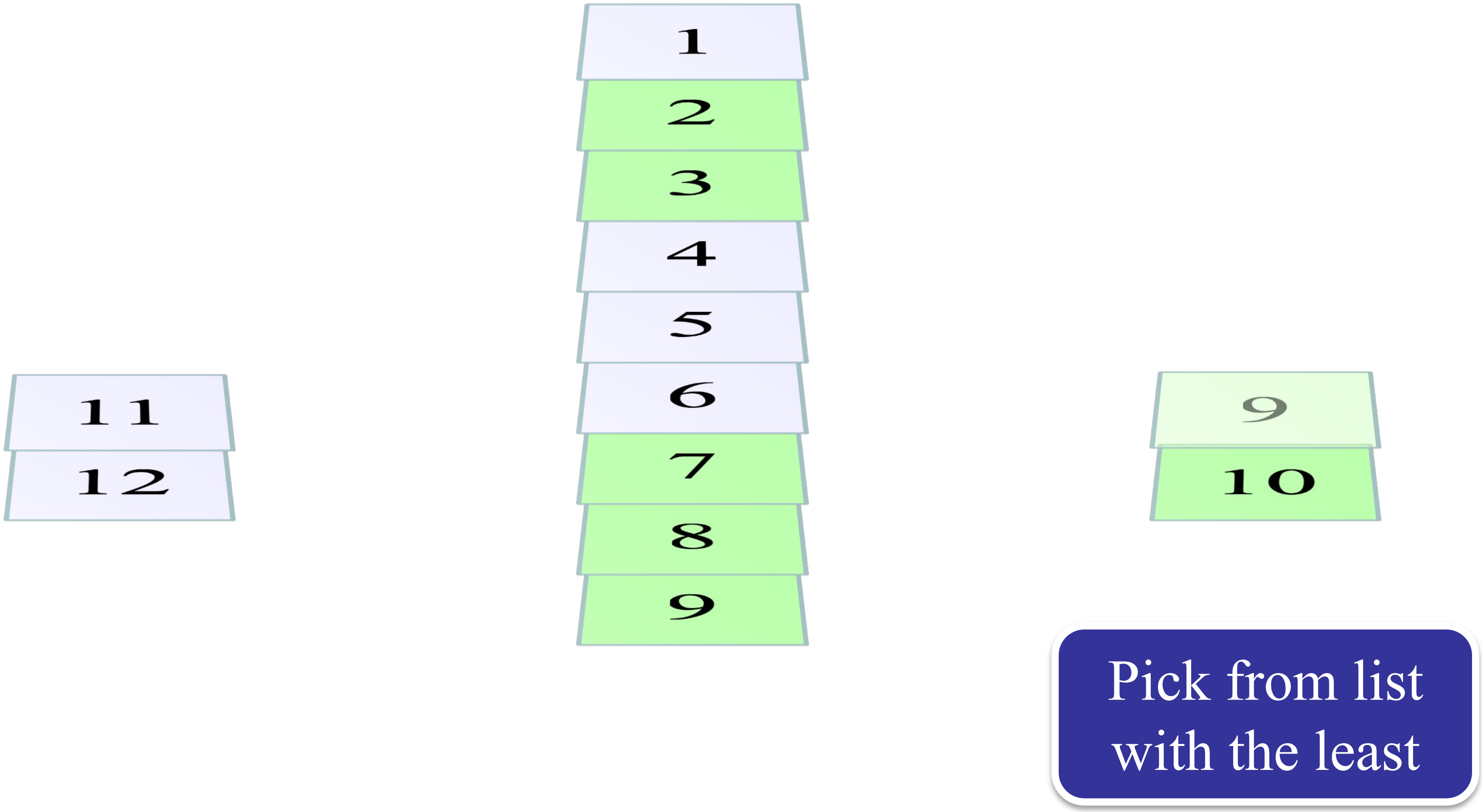

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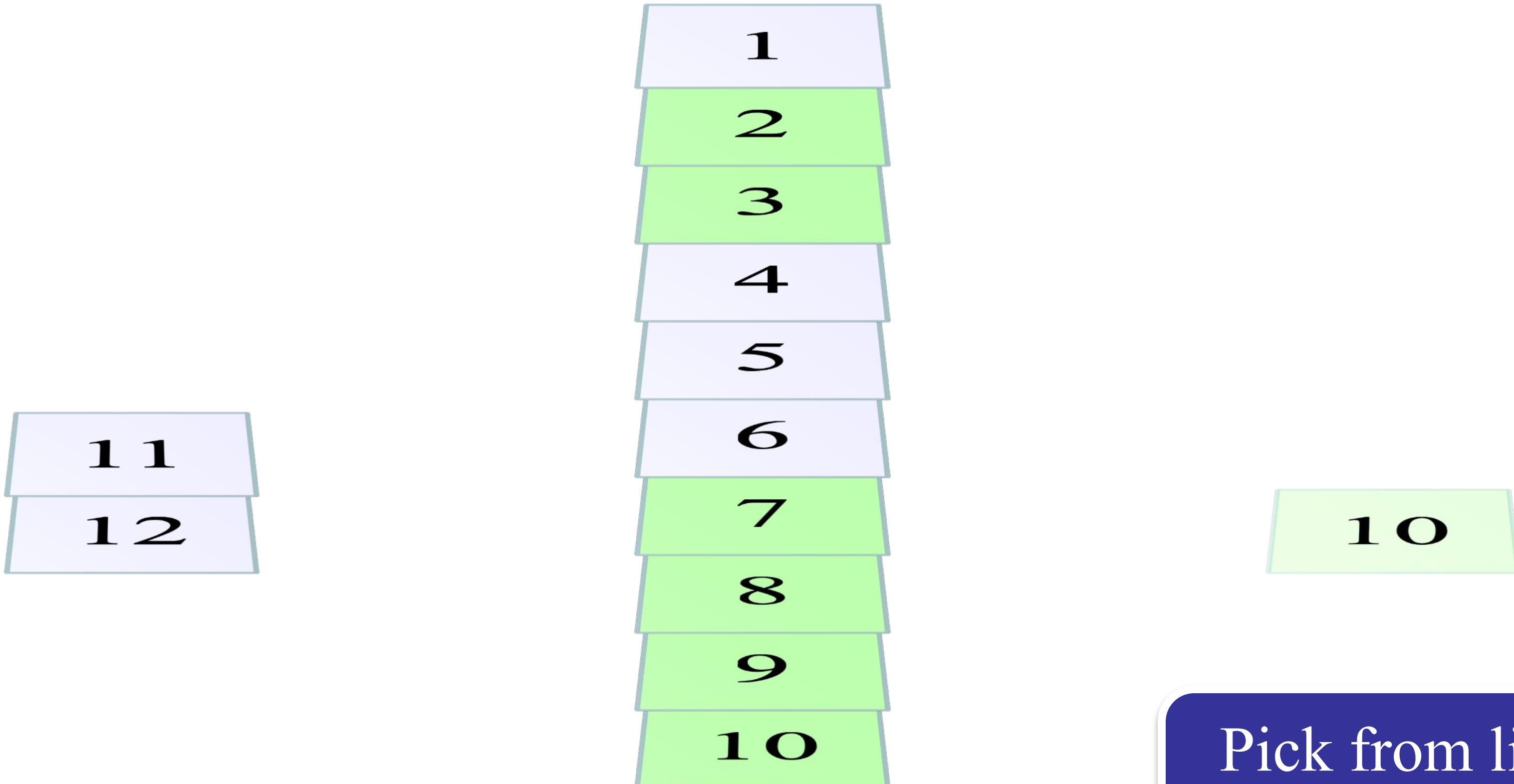

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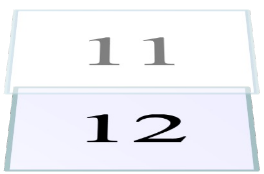
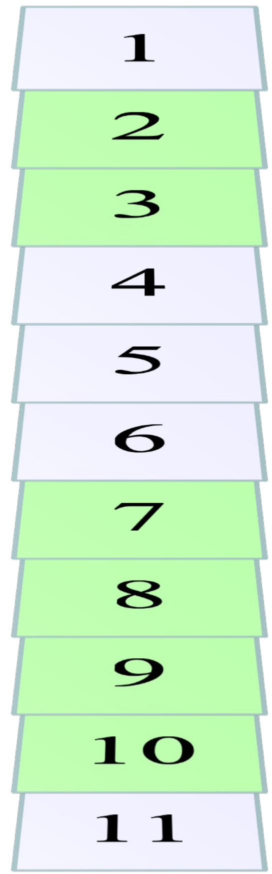
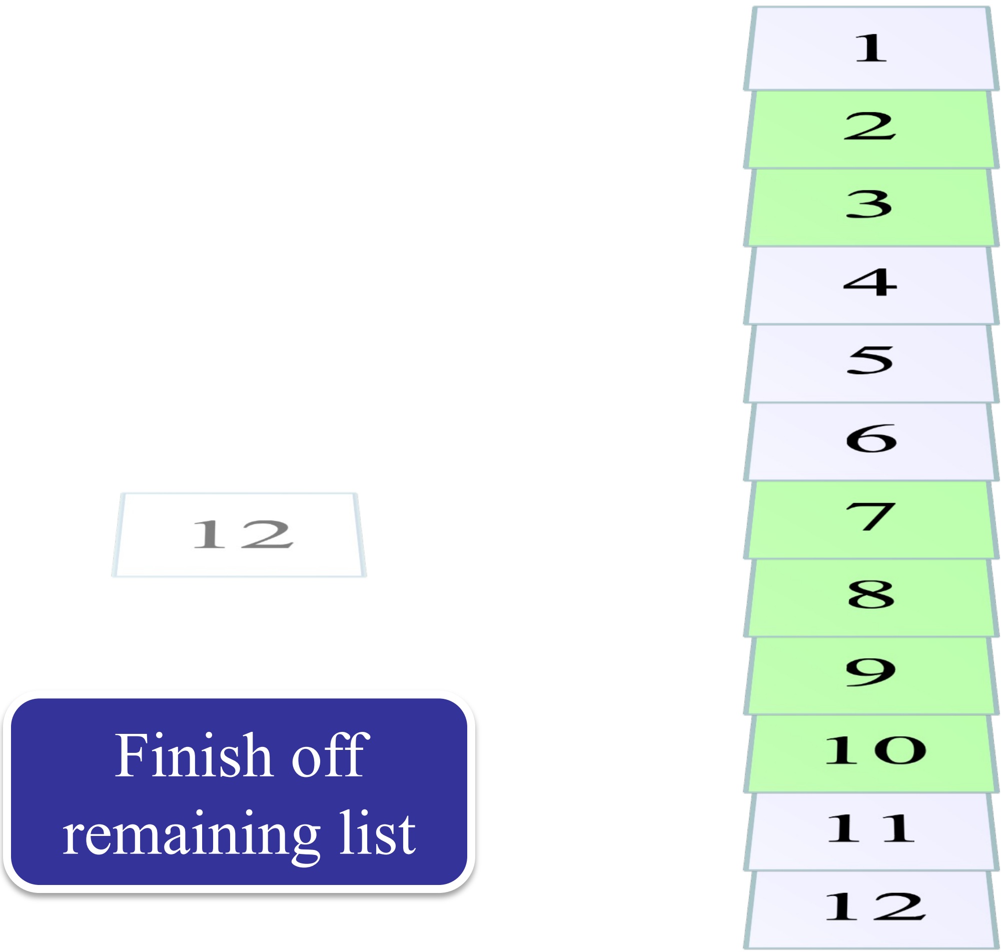


Finish off remaining list



Combining Two Sorted Lists



Merge Sort

```
def merge_sort(b, h, k):
    """Sort the array fragment b[h..k]"""
    if  b[h..k] has fewer than 2 elements:
        return
    # Divide and recurse
    mid = (h+k)//2
    merge_sort (b, h, m)
    merge_sort (b, m+1, k)
    # Combine
    merge(b,h,mid,k) # Merge halves into b
```

- Seems simpler than **qsort**
  - Straight-forward d&c
  - Merge easy to implement
- What is the **catch**?
  - Merge requires a **copy**
  - We did not allow copies
  - Copying takes n steps
  - But so does merge/partition
- n log n **ALWAYS**

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Proof beyond scope of course

What Does Python Use?

- The `sort()` method is **Timsort**
  - Invented by Tim Peters in 2002
  - Combination of insertion sort and merge sort
- Why a combination of the two?
  - Merge sort requires copies of the data
  - Copying pays off for large lists, but not small lists
  - Insertion sort is not that slow on small lists
  - Balancing two properly still gives n log n

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Quicksort is 1959!

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Most of time spent here

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This strategy allows
AI to find even **better**
sorting algorithms